KeyStone Architecture
Telecom Serial Interface Port (TSIP)

User Guide

TEXAS INSTRUMENTS

Release History

Release	Date	Chapter/Topic	Description/Comments
1.0	November 2010	All	Initial Release

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Preface

About This Manual

This manual describes the architecture and functionality of the Telecom Serial Interface Port (TSIP).

Notational Conventions

This document uses the following conventions:

- Commands and keywords are in **boldface** font.
- Arguments for which you supply values are in *italic* font.
- Terminal sessions and information the system displays are in `screen font`.
- Information you must enter is in **`boldface screen font`**.
- Elements in square brackets ([ ]) are optional.

Notes use the following conventions:

Note—Means reader take note. Notes contain helpful suggestions or references to material not covered in the publication.

The information in a caution or a warning is provided for your protection. Please read each caution and warning carefully.

CAUTION—Indicates the possibility of service interruption if precautions are not taken.

WARNING—Indicates the possibility of damage to equipment if precautions are not taken.

Related Documentation from Texas Instruments

C66x CPU and Instruction Set Reference Guide	SPRUGH7
C66x CorePac User Guide	SPRUGW0
TMS320C6000 Programmer's Guide	SPRU198
TMS320C6000 Code Composer Studio Tutorial	SPRU301

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Introduction

1.1 Overview/Purpose of the Peripheral

The TSIP is a multi-link serial interface consisting of a maximum of:

- eight transmit data signals (or links)
- eight receive data signals (or links)
- two frame sync input signals
- two serial clock inputs

Internally, the TSIP offers multiple channels of timeslot data management and multi-channel DMA capability that allow individual timeslots to be selectively processed.

The module can be configured to use the frame sync signals and the serial clocks as redundant sources for all transmit and receive data signals, or one frame sync and serial clock for transmit and the second frame sync and clock for receive. The standard serial data rate for each TSIP transmit and receive data signal is 8.192 Mbps. The standard frame sync is a one or more bit wide pulse that occurs once every 125 µs or a minimum of one serial clock period every 1024 serial clocks.

At the standard rate and default configuration there are eight transmit and eight receive links that are active. Each serial interface link supports up to 128 8-bit timeslots. This corresponds to an H-MVIP or H.110 serial data rate interface. The serial interface clock frequency can be either 16.384 MHz (default) or 8.192 MHz. The data rate for the serial interface links can also be set to 16.384 Mbps or 32.768 Mbps. The maximum number of active serial links is reduced to four and two, respectively, in these configurations. The serial interface clock frequency can be either 32.768 MHz or 16.384 MHz for 16.384 Mbps serial links, and 65.536 MHz or 32.768 MHz for 32.768 Mbps serial links. Maximum occupation of the serial interface links for the entire TSIP is 1024 transmit and receive timeslots in all configurations.

1.2 Features

The TSIP is a serial interface peripheral that connects directly to TEMUX devices with timeslot data management and an integrated DMA capability. The peripheral provides a glueless interface to common telecom serial data streams and efficient internal routing of the data to designated memories in a multi-DSP device. The TSIP provides these functions:

- Direct interface to H-MVIP devices such as TEMUX, ST BUS devices, TSI devices and H.110 compatible devices
- Multichannel transmit and receive of up to 1024 channels
- µ-Law and A-Law companding
- Common and independent framing and clocking for receive and transmit
- Programmable polarity for both frame synchronization and data clocks

1.3 Functional Block Diagram

The TSIP consists of three functional sub-modules as shown in Figure 1-1.
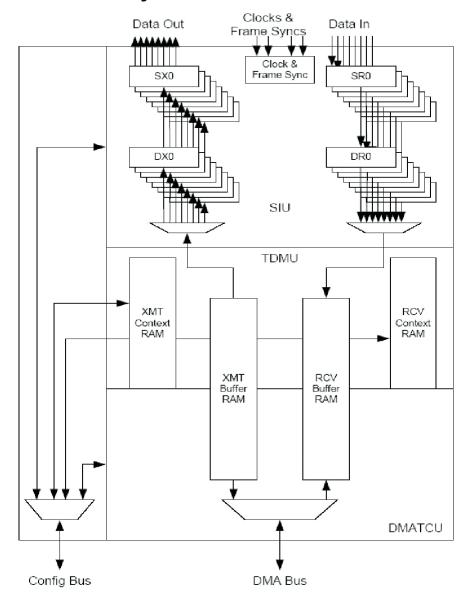
Serial Interface Unit (SIU) — The SIU provides parallel to serial and serial to parallel conversion for transmit and receive respectively.

Timeslot Data Management Unit (TDMU) — The TDMU selectively packs and unpacks timeslot data for receive and transmit based on a channel timeslot definition. There are six transmit and six receive TDMU channels. Each channel is capable of selecting and transferring all of the possible 1024 timeslots in the respective direction.

DMA Transfer Control Unit (DMATCU) — The DMATCU initiates the data transfers between the channel buffers used by the TDMU and the memory buffers used by the DSP.

Figure 1-1   TSIP Block Diagram

1.4 Terminology Used

The following are brief explanations of some terms used in this document:

Term	Meaning
DMA	Direct Memory Access
DMATCU	DMA Transfer Control Unit
SIU	Serial Interface Unit
TDM	Timeslot Data Management
TDMU	Timeslot Data Management Unit
TEMUX	T1-E1 Mux Device
TSIP	Telecom Serial Interface Port

1.5 Memory Map

Table 1-1 lists the summary memory map for the registers for the TSIP.

Table 1-1   TSIP Memory Map

Offset Address	Register Description
0x00000000 - 0x0000007C	TSIP Module Registers
0x00000080 - 0x000000FC	Serial Interface Registers
0x00000100 - 0x0000017C	TDMU Registers
0x00000180 - 0x000001FC	DMATCU Registers
0x00000200 - 0x000003FC	TDMU channel Error Log Registers
0x00000400 - 0x000007FC	Reserved
0x00000800 - 0x00000FFC	TDMU channel Registers
0x00001000 - 0x00001FFC	DMATCU channel Registers
0x00002000 - 0x00007FFC	Reserved
0x00008000 - 0x0000FFFC	TDMU channel Bitmaps
0x00010000 - 0x0001FFFC	TDMU channel Buffers
0x00020000 - 0x0003FFFC	Reserved

Features/Common Architecture

This chapter describes the major features of the TSIP and how it works.

2.1 Interface

The external pin interface of TSIP consists of 20 pins: two clock input pins, two frame sync input pins, eight data input pins, and eight data output pins. The pins are listed in Table 2-1. The module has eight additional control signals to control the output state of the eight data output pins.

Table 2-1 Interface Pins

Pin	Direction	Description
CLK_A	Input	TSIP serial data clock A
CLK_B	Input	TSIP serial data clock B
FS_A	Input	TSIP frame sync A
FS_B	Input	TSIP frame sync B
TR[7:0]	Input	Input TSIP serial data receive. Up to eight serial data streams may be received simultaneously. The serial data clock and frame sync is common for all eight.
TX[7:0]	Output	TSIP serial data transmit. Up to eight serial data streams may be transmitted simultaneously. The serial data clock and frame sync is common for all eight

On the first frame sync, after the TSIP is enabled, the TSIP starts receiving/transmitting on the TR/TX data lane as denoted by the SIU data lanes in Figure 2-1.

Figure 2-1 First Active Frame Sync

2.2 Receive Operation

Data that is received by the TSIP is assembled in a serial to parallel register (DRn) with the most significant bit first, one eight bit register per serial interface input. The data for timeslot 0 is recognized on the first eight serial data clocks following the recognition of the frame sync signal that signals the start of the frame and the programmed delay. The data for timeslot 1 is recognized on the next eight serial data clocks immediately following timeslot 0. Data for successive timeslots are continuously received and recognized on subsequent serial data clocks. The number of data samples processed and buffered depends on the timeslot enable/disable definition in the channel bitmap for each of the receive channels.

2.3 Transmit Operation

Once the TSIP is configured for transmission, on the first frame sync after the TSIP is enabled, its DMA engine reads the data from the memory buffers and stores it in internal channel buffers. It starts transmitting from the very next frame sync, depending on the timeslots enabled in all the six channels. The data for timeslot 0 is sent on the first eight serial data clocks following the recognition of the frame sync signal that signals the start of the frame and the programmed delay. Data that is transmitted by the TSIP is output from a parallel to serial register with the most significant bit first, one 8-bit register per serial interface output. The data for timeslot 1 is sent on the next eight serial data clocks immediately following timeslot 0. Data for successive timeslots are continuously transmitted on subsequent serial data clocks. The

eight transmit data registers, one for each serial interface signal, mean that the TSIP executes its load operation up to eight times for each timeslot. The number of actual data samples loaded depends on the timeslot enable/disable definition in the channel bitmap for each of the channels. With the data for each timeslot, an enable/disable indication is also provided. The enable/disable indication is used by the serial interface to control the state of the output. Three output states for disabled timeslots are supported: driven high, driven low, and high impedance. The disabled timeslot output state selection is defined by XMTDIS as shown in Table 2-2.

Table 2-2 Transmit Disabled State

Bit	Field	Value	Description
9-8	XMTDIS([1])		Transmit output disable state
		00	High impedance
		01	Reserved
		10	Driven low
		11	Driven high

1. This field is write protected when SIU transmit is enabled

When an enabled timeslot immediately follows a disabled timeslot it is sometimes desirable to delay the enable of the output buffer for the first bit transmitted on the enabled timeslot. The delay is defined by XMTDLY as shown in Table 2-3. The selection determines whether the output buffer is turn on immediately for the first bit of the enabled timeslot or whether it is turned on after a delay of one-half serial clock period.

Table 2-3 Enable Output Delay Selection (XCR)

Bit	Field	Value	Description
10	XMTDLY([1])		Transmit output enable delay. Defines fixed delay for enabling the output buffer on the first bit of an enabled timeslot that immediately follows a disabled timeslot. This control applies only when the disable state for the output buffer is high impedance.
		0	No added delay
		1	Output enable is delayed by one-half serial clock period.

1. This field is write protected when SIU transmit is enabled.
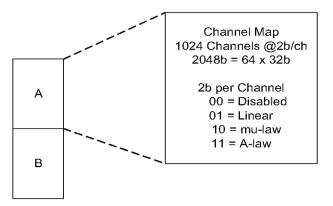
2.4 Bitmap Context

The receive operation of the TSIP reads and processes the data in the DRn registers and stores the results in the channel buffer based on the channel bitmap context. The registers are read sequentially with the sequence order defined in the SIU by the selected data rate option. Data received on each of the receive serial links can be stored in none, some, or all of the channel buffers. On reaching a watermark in the channel buffer, its DMA engine transfers the data from the channel buffers to the memory buffers outside the system. Each transmit/receive channel has dual buffers, PING and PONG, that are shared with the DMATCU. Each of the buffers is managed as a FIFO memory and consists of 64 32-bit words of memory. During the processing of a given frame the TSIP processes the data from one of the FIFO memories based on whether the PING or PONG buffer is active. The channel buffers are not software accessible once the TDMU has been enabled.

The channel bitmap consists of 64 32-bit words of memory for a total of 2048 bits. Each pair of bits controls the enable/disable state of a physical timeslot. When a given di-bit value is 00b the corresponding physical timeslot is disabled for that channel. Any other di-bit value represents an enabled channel: 01b indicates that the data is 8-bit linear, 10b indicates that the data should be μ-law expanded, and 11b indicates that the data should be A-law expanded.

Figure 2-2   Channel Bitmap

For a given physical timeslot the corresponding di-bits are read from the channel bitmap for each of the channels. If the physical timeslot is enabled for the channel the data, either 8 or 16 bits, is read from the FIFO memory. If 16-bit data is read from the FIFO memory the data is compressed to 8 bits using either μ-law or A-law compression depending on the enable selection. At transmit side, data samples from multiple timeslots can be loaded from the channel buffer (FIFO memory) into a channel register and unpacked for enabled timeslots. The resulting 8-bit sample data for each of the enabled channels is then OR'd together to create a single 8-bit sample value that is written to the appropriate transmit data register (DXn). The ability to combine the data from multiple channels enables the creation and use of physical timeslot sharing provided that the bit values are zero for the unused bits in the physical timeslot data for a given channel. Similarly at the receive side, if the same timeslot is enabled on multiple channels, data from receive register is loaded into the channel buffers of all the channels for which the timeslot is enabled. Any 8-bit data is read/written from/to the next byte location in the FIFO memory. Any 16-bit data is read/written from/to the next 16-bit word location in the FIFO memory.

2.4.1 A-law and μ-law Companding

When receive data is μ-law or A-law expanded it is stored left-justified in a 16-bit word. The two or three least significant bits are always zero for μ-law and A-law, respectively. Likewise, transmit data that is to be μ-law or A-law compressed should be left justified. The two or three least significant bits are truncated for the compression. Expansion and compression is completed based on the G.711 standard including the channel coding (bit inversions for ones density).

Since the G.711 standard does not distinguish the μ-law expansion of 0xFF and 0x7F, a special case is created in the TSIP. Both normally result in a 16-bit linear value of 0x0000. For data that undergoes signal processing this is not a problem. However, for data that is compressed and looped back the compression results in 0xFF and the original 0x7F code is lost wherever it existed. To prevent this problem, the expansion of 0x7F is redefined to be 0xFFFF and the algorithmic compression of the 16-bit value 0xFFFF, which normally results in 0x7E, is treated as a special case to allow 0x7F to be produced. This does not result in any loss in the coding based on the G.711 standard since the G.711 standard only defines the data conversion for 14-bit linear data values and the two least significant bits of the left-justified value in a 16-bit linear format otherwise would be truncated.

2.5 Serial Link Processing and Buffering

When the selected data rate option is one of the 8 Mbps options, all eight serial interface inputs are active. Links refer to the number of transmit and receive I/Os supported by the TSIP module. Active links refer to the number of transmit and receive I/Os that are usable at a given data rate. Figure 2-3 illustrates the channel bitmap definition for both 8 Mbps data rate options. The nomenclature defines the logical timeslot, LTSxxx, and physical timeslot, Ln-xxx, where n represents the link number and xxx represents the physical timeslot. When the selected data rate option is the 16 Mbps option four serial interface inputs are active, but data is read from all of the DRn/Dxn registers.

Figure 2-3   8 Mbps Channel Definition (8 Links with 8 Active Links)
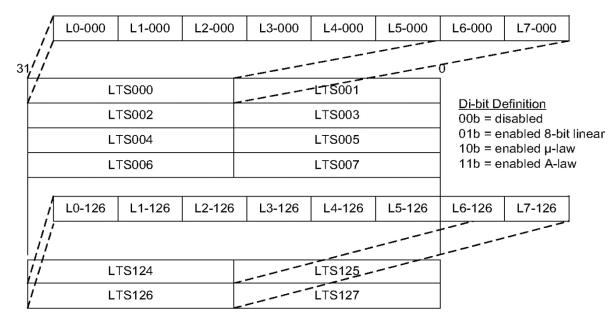
Figure 2-4 illustrates the channel bitmap definition for the 16 Mbps data rate. When the selected data rate option is the 32 Mbps option two serial interface inputs are active. Once again, data is read from all of the DRn/Dxn registers. Figure 2-5 illustrates the channel bitmap definition for the 32 Mbps data rate. Data samples from multiple timeslots can be merged in a channel register prior to actually writing to the channel buffer (FIFO memory) when a 32-bit word is fully packed.

Figure 2-4   16 Mbps Channel Definition (8 Links with 4 Active Links)
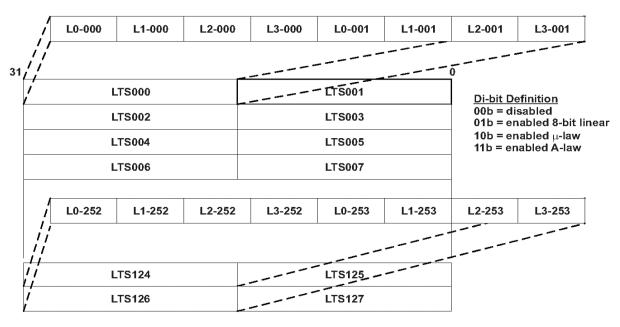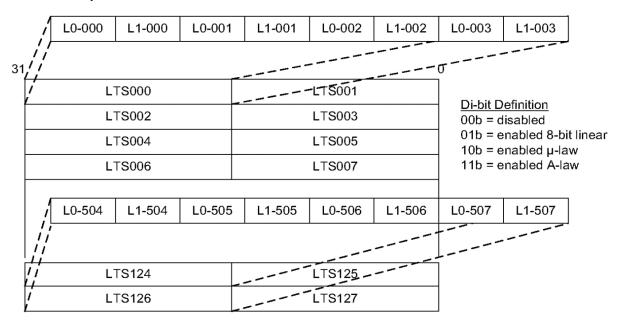
Figure 2-5   32 Mbps Channel Definition (8 Links with 2 Active Links)

2.6 Multi-Channel Operation
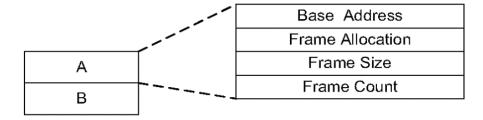
The timeslot data management unit (TDMU) of the TSIP is a multi-channel unit that manages the data for the SIU. There are six transmit and receive channels in a total of twelve channels. For each transmit channel the timeslot context is defined to support enable/disable control of all timeslots on all of the serial outputs operating at an 8.192 Mbps data rate or 128 8-bit timeslots. For each receive channel the timeslot context control is defined to support enable/disable control of all timeslots on all of the serial inputs operating at an 8.192 Mbps data rate or 128 8-bit timeslots.

Each DMA channel has dual buffers, PING and PONG, that are shared with the TDMU. The use of the dual buffers allows memory buffers, including each frame of data within the memory buffer, to begin at designated locations in memory. When the next frame sync event occurs the DMATCU flushes the remaining contents of the current channel buffer and then switches focus to the second channel buffer. The channel buffers are not software accessible once the DMATCU has been enabled.

2.6.1 Transmit and Receive Channel Context

The context control for each channel consists of a dual set of configuration registers (A and B). These sets of four configuration registers each become active in direct relationship to the active channel bitmap of the TDMU. The registers define the memory address for the data buffers (Base Address), frame size allocations (Frame Allocation and Frame Size), and number of frames allocated (Frame Count) within the memory buffers, as shown in Figure 2-6.

Figure 2-6   Channel Configuration

Each base address register is a 32-bit register representing the base address of a frame buffer in memory. Since data is transferred from the TSIP to memory in 32-bit elements and addresses correspond to byte addresses, the two least significant address bits of the base address must always be zero. The frame buffer in memory consists of storage for multiple frames of data. The number of live frames that can be stored in the frame buffer is defined by the frame count. Frames of data are stored in the frame buffer with the first frame beginning at the base address. When a context switch occurs, the storage based on the new context always begins at the base address defined by that context.

The next frame is stored in the frame buffer beginning at the location defined by the base address plus the frame allocation. When the number of frames defined by the frame count have been stored in memory, the next frame is stored beginning at the base address. Since data transfers are always in 32-bit elements and the frame allocation is an address offset to the base address, the two least significant bits of the frame allocation must be zero.

The frame size is not used explicitly in receive channel transfers. It is present for context consistency and is described with the transmit channel context. With respect to the context, the frame size represents the amount of storage a single frame of data actually occupies, including the 32-bit CID/FRFC synchronization values at the beginning and end of the frame. Realistic frame size values begin at twelve when there is at least one timeslot enabled. The minimum acceptable value for frame size is eight, which implies that there are no real data samples for the frame since eight bytes are required for the two 32-bit synchronization values. A frame size that is less than eight is logged as an error for a channel. In the event that the frame size exceeds the frame allocation there is a potential overlap in the frame data of two frames. The DMATCU detects this condition and logs an error. The DMATCU continues to function using the defined values.

A channel active status register consisting of pairs of bits for each channel (CHn) identifies which channel configuration registers (A or B) is active at any given time for each channel. When a channel configuration (A or B) is active those configuration registers are write protected. A channel configuration (A or B) that is pending active also is write protected. Write accesses to the inactive configuration registers update the inactive registers. Read accesses from either the active or inactive configuration complete as normal. The status of each CHn field is defined in Table 2-4.

Table 2-4   Channel Configuration Active Status

Bit	Field	Value	Description
	CHnCS		Transmit/Receive channel n active status. Defines whether A and B configurations are in active status.
		00	Both A and B configurations are in active status.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved

When the channel is disabled both bitmaps are inactive and can be modified. The least significant bit of the CnID value determines which of the two channel bitmaps is made pending and becomes active on the next superframe. A channel bitmap can be designated as pending active at any time. When a bitmap becomes pending active, both it and the active bitmap are protected until the pending active bitmap is made active and the active bitmap is made inactive. Changes in the actual active assignment always occur on a superframe. Writes to the channel enable register to set the A or B bitmap as pending active must have the CHnENB set to enable the channel or to keep the channel enabled if it is already running. Writes to the channel enable register with a new CnID value in which the least significant bit is identical to the previous CnID value causes the CID value used beginning with the next superframe to be the new value. However, the channel bitmap and DMA context that are currently active remain unchanged. Writes to the channel enable register that clear the channel enable bit take effect on the next frame event.

Table 2-5   Channel Enable Register

Bit	Field	Value	Description
15-8	CnID	0x00	Configuration identification. Defines the configuration selection that takes effect on the next superframe synchronization event. The least significant bit determines whether the A or B channel bitmap and configuration registers are used.
0	CHnENB		Channel n enable
		0	Channel n is disabled/stopped.
		1	Channel n is enabled/running.

2.6.2 Frame Format

The first 32-bit location in the frame in the channel buffer is filled with a 32-bit synchronization value. This 32-bit value consists of an 8-bit channel configuration ID (CnID) and the lower 24 bits from the free running frame counter. The channel configuration ID is defined by the software in the corresponding channel enable register. At the end of a receive frame, the beginning of the next frame, when the channel buffer is flushed to memory the same 32-bit value is written once again to the channel buffer. At the same time, a new 32-bit value with at least the incremented free running frame counter is written to the other half of the channel buffer. The presence of the synchronization value at the beginning of the frame in memory provides a way for software to recognize and track the configuration that was used for that frame. The presence of the same value at the end of the frame provides a way for software to confirm that the entire frame has been written to memory.

Figure 2-7, Figure 2-8, and Figure 2-9 illustrate the way the data packing takes place within the channel buffer based on the enabled channels. Figure 2-7 is a list of enabled physical timeslots for a single receive channel. In this example, the receive data on serial link 2 (TR2) is to be expanded using either m-law or A-law expansion. The receive data from all other serial links is to be stored as 8-bit linear. The first 32-bit location in the channel buffer is filled with a 32-bit synchronization value.

The data for the first eleven enabled timeslots that are received is 8-bit linear and the data is packed in the channel buffer on successive bytes. As the example illustrates in Figure 2-8 and Figure 2-9, when timeslot 5 for TR2 (L2-005) is received, expanded, and stored the next available location is an odd byte location so a dummy byte is inserted and L2-005 is stored on a 16-bit boundary. When L2-045 and L2-046 are received the next available location is already 16-bit aligned resulting in no extra dummy bytes. However, when L2-047 arrives the next available location is once again an odd byte resulting in the insertion of a dummy byte to allow L2-047 to be stored on a 16-bit boundary. At the end of a frame pad bytes are used to align the final 32-bit synchronization value. The actual content of the dummy bytes and pad bytes is undefined.

Figure 2-7   Example Timeslot List by the TSIP Serial Link for 1 Channel

Link 0	Link 1	Link 2	Link 3	Link 4	Link 5	Link 6	Link 7
004	004		004			004	
005			005	005			
	008				008		
009	009	009	009	009	009		
	024						
03							
	040		040				
				043			
		045					
		046			046	046	
047		047					
	052		052	052		052	
101							
			102				

Figure 2-8   Channel Buffer (BEND=1)

31			0
L0-004	L1-004	L3-004	L6-004
L0-005	L3-005	L4-005	L1-008
L5-008	L0-009	L1-009	dummy
L2-009		L3-009	L4-009
L5-009	L1-024	L0-039	L1-040
L3-040	L4-043	L2-045	
L2-046		L5-046	L6-046
L0-047	dummy	L2-047	
L1-052	L3-052	L4-052	L6-052
L0-101	L3-102	pad	pad
CID	FRRC		

Figure 2-9   Channel Buffer (BEND=0)

31			0
CID	FRRC		
L6-004	L3-004	L1-004	L0-004
L1-008	L4-005	L3-005	L0-005
dummy	L1-009	L0-009	L5-008
L4-009	L3-009	L2-009	
L1-040	L0-039	L1-024	L5-009
L2-045		L4-043	L3-040
L6-046	L5-046	L2-046	
L2-047		dummy	L0-047
L6-052	L4-052	L3-052	L1-052
pad	pad	L3-102	L0-101
CID	FRRC		

2.7 Interrupts

Two types of interrupts are generated by the TSIP. The first is an indication of data transfer progress. The second is an indication of error conditions within the module or due to common bus architecture (CBA) status codes.

Each DMA channel has a set of these interrupts that are routed to its corresponding core. Refer to device specific datasheet for detail.

2.7.1 Frame Interrupts

Each time the timeslot counter rolls over for the transmit serial links or the receive serial links the corresponding frame interrupt is made pending. For receive, some (all) of the data for the frame that was just completed prior to the frame sync is in the channel buffer. At least one additional transfer needs to be made to flush the data from the buffer. An acknowledgment of the transfer completion, assuring that the data has actually been written to memory is returned to the TSIP. The TSIP offers the option of asserting the frame interrupt that has been made pending as a result of the frame sync upon receipt of acknowledge for the final transfer, shown in Table 2-6. A programmable timeslot delay, shown in Table 2-7, is also available to determine when the interrupt is actually asserted. Based on the telecom interface timing references this offers the option of generating an interrupt every 125 μs.

Table 2-6   Frame Interrupt Control (RTDR and XTDR)

Bit	Field	Value	Description
13-12	RCVFRINT[1]		Receive frame interrupt selection
		00	Interrupt asserted on transfer acknowledge only.
		01	Interrupt asserted on delay count only.
		10	Interrupt asserted on either transfer acknowledge or delay count.
		11	Interrupt asserted on both transfer acknowledge and delay count.
13-12	XMTFRINT[1]		Transmit frame interrupt selection
		00	Interrupt asserted on transfer acknowledge only.
		01	Interrupt asserted on delay count only.
		10	Interrupt asserted on either transfer acknowledge or delay count.
		11	Interrupt asserted on both transfer acknowledge and delay count.

1. This field is write protected once DMA is enabled.

Table 2-7   Timeslot Delay Counter (RTDR and XTDR)

Bit	Field	Value	Description
6-0	RCVFDLY[1]	0x07F	Receive frame interrupt delay [0-127]
			n + 1 Number of logical timeslots the frame interrupt is delayed after going pending.
6-0	XMTFDLY[1]	0x07F	Transmit frame interrupt delay
			n + 1 Number of logical timeslots the frame interrupt is delayed after going pending.

1. This field is write protected once DMA is enabled.

2.7.2 Superframe Interrupt

The frequency of this interrupt is independently programmable with the frame sync counter. The superframe interrupt is made pending on the frame sync counter rollover. The same programmable timeslot delay used with the frame interrupt determines when the interrupt is actually asserted. The actual assertion of the superframe interrupt can be determined solely by the delay, by acknowledgment that the final data transfer for the last frame in the superframe is complete or a combination of both as shown in Table 2-8.

Table 2-8   Superframe Interrupt Control (RTDR and XTDR)

Bit	Field	Value	Description
15-14	RCVSFINT[1]		Receive superframe interrupt selection
		00	Interrupt asserted on transfer acknowledge only.
		01	Interrupt asserted on delay count only.
		10	Interrupt asserted on either transfer acknowledge or delay count.
		11	Interrupt asserted on both transfer acknowledge and delay count.
15-14	XMTSFINT[1]		Transmit superframe interrupt selection
		00	Interrupt asserted on transfer acknowledge only.
		01	Interrupt asserted on delay count only.
		10	Interrupt asserted on either transfer acknowledge or delay count.
		11	Interrupt asserted on both transfer acknowledge and delay count.

1. This field is write protected once DMA is enabled.

2.7.3 Error Interrupt

Error interrupts are asserted when certain fault conditions are identified. There is one error interrupt per channel. An error interrupt queue for each channel provides a history of up to sixteen error interrupts per channel. Each item in the queue includes an 8-bit error code and 24 bits of additional information. Error codes are listed in Table 2-9.

2.7.4 Exception Conditions

Table 2-9   Error Codes

Error Code	Info Field	Category	Description
0x00			No Error
0x01	RCVFRFC	1	Receive Overrun at the beginning of frame
0x02	RCVFRFC	1	Receive Overrun during the frame
0x03	CIDA	3	Receive Overrun-Enable channel bitmap greater than Size
0x05	CIDA	3	Receive Size less than 8
0x10	CIDA & CIDF	2	Transmit configuration ID mismatch
0x11	XMTFRFC	1	Transmit Underrun at the beginning of frame
0x12	XMTFRFC	1	Transmit Underrun during the frame
0x13	CIDA	3	Transmit Underrun-Enabled channel bitmap greater than Size
0x15	CIDA	3	Transmit Size less than 8
0xF1-0xF7			CBA read status (rstatus) errors. (The three least significant bits of this error code are defined in the CBA 3.0 Specification.)
0xF9-0xFF			CBA write status (sstatus) errors. (The three least significant bits of this error code are defined in the CBA 3.0 Specification.)

2.7.4.1 Receive Errors

Error interrupts are generated on DMATCU and TDMU receive channels due to buffer overrun and context mismatch conditions. A buffer overrun occurs anytime there is no space in the channel buffer for an enabled timeslot.

- **Error Code = 0x01**. This code indicates that the channel buffer is not empty at the beginning of the frame. It is the most serious overrun condition because it indicates that there is receive data in both the PING and PONG buffers. The additional information is the current value of the receive free running frame counter RCVFRFC).

- **Error Code = 0x02**. This code indicates that the channel buffer is full during a frame transfer and the TDMU is unable to add data for an enabled receive timeslot. The additional information is the current value in the receive free running frame counter (RCVFRFC).

- **Error Code = 0x03**. This code indicates that the number of enabled timeslots in the channel bitmap results in a buffer requirement that exceeds the frame allocation or frame size. The additional information is the current (active) channel configuration ID (CIDA).

- **Error Code = 0x05**. This code indicates that the frame size is less than 8, which is the minimum for a channel with no timeslots enabled. The additional information is the current (active) channel configuration ID (CIDA).

2.7.4.2 Transmit Errors

Error interrupts are generated on transmit channels due to buffer underrun and context mismatch conditions. A buffer underrun occurs anytime there is no data in the channel buffer for an enabled timeslot.

- **Error Code = 0x10**. This code indicates that the configuration identity in the transferred frame data does not match the configuration identity in the active context. The additional information is the current (active) channel configuration ID (CIDA) and the transferred frame data channel configuration ID (CIDF).

- **Error Code = 0x11**. This code indicates that the channel buffer is empty at the beginning of the frame. It is the most serious underrun condition because it indicates that both the PING and PONG buffers are empty. The additional information is the current value in the transmit free running frame counter (XMTFRFC).

- **Error Code = 0x12**. This code indicates that the channel buffer is empty during a frame transfer and the TDMU is unable to provide data for an enabled transmit timeslot. The additional information is the current value in the transmit free running frame counter (XMTFRFC).

- **Error Code = 0x13**. This code indicates that the number of enabled timeslots in the channel bitmap results in a buffer requirement that exceeds the frame allocation or frame size. The additional information is the current (active) channel configuration ID (CIDA).

- **Error Code = 0x15**. This code indicates that the frame size is less than 8, which is the minimum for a channel with no timeslots enabled. The additional information is the current (active) channel configuration ID (CIDA).

2.7.4.3 CBA Errors

When an error occurs on a CBA transaction a 3-bit status code is returned by the slave to the master. The master for the TSIP is the DMA interface, which is driven by the DMATCU. The status codes are provided for read and write transactions independently. A status code of 0 is an indication of success for both reads and writes and no entry is made in the error interrupt queue. When any other status code is returned the code is added to 0xF0 for reads and 0xF8 for writes and stored in the error interrupt queue.

2.7.5 Error Queue Management

The error queue for each channel is maintained in FIFO format with a write to pop requirement. The write to pop an entry from the queue allows the top of the queue to be inspected in a watch window during an emulation halt without removing the entry from the queue. For each queue two memory mapped registers are used to support queries of the queue. The first register is a control or queue management register. The second register includes a count of the number of errors on a given queue. A third register location is used to represent the top of the error queue, or the oldest error on the queue. When an error occurs the error code and information are added to the queue for the appropriate channel and the error interrupt is asserted. If the queue is full the new error is discarded and an overflow is indicated. Once the overflow indication is set all subsequent errors are discarded, even if there is space on the queue, until the DSP (user) clears the overflow condition. When a DSP responds to an error interrupt it performs the following steps.

1. The count register for the channel error queue is read and the number of error entries on the queue is obtained.
2. The error queue register is read to obtain the top (oldest) error entry on the queue.
3. The entry is removed or popped off the queue by doing a write to the error queue register.
4. If the count read in step 1 is greater than 1 then steps 2 and 3 should be repeated until all error entries have been read from the queue.

Reading the error queue register when the count register is zero should return an error entry value of zero.

2.7.6 Error Recovery

In the case where a frame sync pulse occurs before the frame period (1024 serial data clocks for the default configuration) has completed, that frame sync pulse is ignored. The 1024 serial data clocks in the default configuration represent 128 8-bit timeslots. Once all of the required timeslots have been counted, the next frame sync pulse defines the beginning of a new frame. In the case where a frame sync pulse does not occur until some number of clock periods after the frame period has completed, the additional serial data clocks are ignored.

During the time when the serial data clocks are ignored, no timeslot events are created and no data is clocked in or out of the SIU. When the frame sync pulse does occur, the bit counter begins counting from its initial value with the appropriate delay to the frame sync input.

In the case where the serial data clock is removed or aborted in some other way, the data for at most one frame is corrupted although there may be one or more frame periods of missing data. When the serial data clock is restored, the number of clock periods needed to complete the frame that was in progress when the serial data clock was removed are counted. Any data received and transmitted during this period is considered corrupted.

The elapsed time also corresponds to additional missing frames of data. Any frame sync pulses that occur during that time are ignored. Once enough serial data clocks have occurred to correspond to the frame period, the serial interface recognizes the next frame sync input and data is received and transmitted starting with that frame sync.

2.8 Endianness

The way in which the time ordered data is packed and unpacked in the channel buffer is determined by the endian mode selection (BEND), as shown in Table 2-10. The 32-bit synchronization value that is also included in the channel buffer is not affected by the endian mode selection. Figure 2-8 and Figure 2-9 show the resulting packed data in the channel buffer for the two endian modes. The 32-bit synchronization value (CID+FRFC) is stored in the same manner independent of the endian selection. In addition, the 16-bit expanded data samples are stored in the same manner within a 16-bit field independent of the endian selection. There is no byte swapping within the designated fields for either 32-bit or 16-bit values.

Table 2-10   Endian Mode Selection

Bit	Field	Value	Description
0	BEND([1])		Endian mode selection. Bend = iso bend + 1.
		0	Little endian format
		1	Big endian format

1. This field is write protected once DMA is enabled.

2.9 Priority Control

More than one channel can have its transfer request flag set. When there are multiple transfer request flags set, round-robin arbitration is used to select the next channel for transfer. To sequence all transfers (reads and writes) the winning receive channel may also need to win arbitration with a transmit channel. Priority for contested receive and transmit channel requests within the TSIP is traded back and forth with each contested request. The priority of the transfer request made by the TSIP to access memory is defined by PRI as shown in Table 2-11.

The priority is escalated by one additional level, if possible, each time the channel buffer fills to the point that another transfer request could be generated while an existing request remains outstanding. Escalation is limited by MAXPRI, also shown in Table 2-11. MAXPRI is expected to be a value that is less than or equal to PRI. When MAXPRI and PRI are equal no escalation is made. If MAXPRI is greater than PRI, the programmed value is ignored and MAXPRI is treated as though it is equal to PRI.

When the last transfer request has been generated (that is, when the completion of the transfer request leaves fewer than eight 32-bit words in the channel buffer) the normal priority, defined by PRI, is restored. For example, a transfer request is made at the normal priority when the channel buffer initially contains eight 32-bit words to be transferred to memory. In the event that an additional eight 32-bit words are added to the same channel buffer before the transfer is completed for the first request, the priority is escalated one additional level. If there are fewer than eight 32-bit words remaining in the channel buffer when the second request has been generated, the priority is reduced back to the normal priority level for the next request.

Table 2-11   DMATCU Transfer Priority (DMATCU_CGR)

Bit	Field	Value	Description
6-4	MAXPRI[1]		Maximum transfer priority
		0	Priority 0-Highest priority
		1	Priority 1
		...	Priority 2-6
		7	Priority 7-Lowest priority
2-0	PRI[1]		Transfer priority
		0	Priority 0-Highest priority
		1	Priority 1
		...	Priority 2-6
		7	Priority 7-Lowest priority

1. This field is write protected once DMA is enabled

Clocks, Frames, and Data

This chapter procides details on frame and clock operation, data rates, data delay, and synchronization.

3.1 Frame and Clock Operation

The clock is expected to be a continuously running clock with uniform clock periods and with exactly the timeslot count times the bit count number of clock periods between frame sync pulses. With the default counter values that are used with 8.192 Mbps links this period is exactly 1024 clocks. The TSIP anticipates the frame sync pulses to come at that interval. When the frame sync pulse occurs early or late the response of the TSIP is defined to assure that the operation recovers as gracefully as possible.

3.2 Data Rate Options

The TSIP offers six data rate options. The data rate options are listed in Table 3-1 with some corresponding information. Data rate options are available to support double clocking (default) or single clocking of the serial data for 8.192, 16.384, or 32.768 Mbps serial links. When the serial data rate changes the maximum number of active serial data links is also changed. There is an inverse relationship between the two: 8 active links at 8.192 Mbps, 4 active links at 16.384 Mbps, or 2 active links at 32.768 Mbps. However, the total number of 8-bit timeslots is constant for all data rate options. When the number of active links is reduced the output buffers for the inactive links are put in the defined disabled state. The defined disabled state for inactive links is the same as the one defined for disabled timeslots.

Table 3-1   Data Rate Options

Option	Name	Active Links	Data Rate	Clock Frequency	Timeslots/Link
000	8m16	8	8.192 Mbps	16.384 MHz	128
001	8M8	8	8.192 Mbps	8.192 MHz	128
010	16M32	4	16.384 Mbps	32.768 MHz	256
011	16M16	4	16.384 Mbps	16.384 MHz	256
100	32M64	2	32.768 Mbps	65.536 MHz	512
101	32M32	2	32.768 Mbps	32.768 MHz	512

3.3 Clock and Frame Sync Redundancy

The receive and transmit interfaces share a clock, frame sync, and data rate by default. This setup offers the option of redundant (primary and secondary) clock and frame sync sources that can be switched when the primary source fails. An option allows the receive and transmit interfaces to be configured independently with each having its own clock and frame sync and possibly data rates.

The default configuration uses one clock and frame sync as the primary and the second clock and frame sync as the secondary. This configuration uses the primary clock and frame sync to control the transmit data on all eight data outputs and the receive data on all eight data inputs. The use of a primary and secondary serial clock and frame sync allows the serial links to be maintained with a small impact on the data. The impact is limited to the loss of a few frames of data when a switchover is required.

Transitions between the use of the primary clock and frame sync and the backup clock and frame sync are handled by software changing the configuration based on a message it receives. The selection of the clock and frame sync also can be independent for the transmit serial links and the receive serial links. This provides a second configuration option in which the clock and frame sync used for transmit is different from the clock and frame sync used for receive. In this configuration there is no secondary or backup source.

The selection of redundant or independent clocking is made using CLKD as shown in Table 3-2.

Table 3-2   Clock Redundancy Selection (SIU_GCR)

Bit	Field	Value	Description
4	CLKD		Clock redundancy selection.
		0	Redundant clock mode. Transmit and receive share the same clock and frame sync. The common selections are those defined by XMTSRC, SMTCLKM, and MTDATR in the transmit select and enable register. The RCVSRC, RCVCLKM, and RCVDATR fields are ignored.
		1	Dual clock mode. Transmit and receive are specified totally independently.

3.4 Single and Double Data Rate

Data rate options are independent for receive (RCVCLKM and RCVDATR) and transmit (XMTCLKM and XMTDATR) when CLKD is 1. RCVCLKM and XMTCLKM define whether a double-rate or single-rate clock is being provided for receive and transmit, respectively.

For a double rate clock there are two sample clock edges per data bit. Data is sampled on any of the two edges, depending on the edge that frame sync is detected and also on the polarities of the data sampling clock edge configuration.

RCVDATR and XMTDATR define the actual data rate (or timeslots per link) for receive and transmit as shown in Table 3-3.

Table 3-3   Single Rate and Double Rate Clock Options

Bit	Field	Value	Description
0	RCVCLKM		Receive clock mode
		0	TCLK_x is a double rate clock.
		1	TCLK_x is a single rate clock
0	XMTCLKM		Transmit clock mode
		0	TCLK_x is a double rate clock.
		1	TCLK_x is a single rate clock
2-1	RCVDATR		Receive data rate
		00	8.192 Mbps
		01	16.384 Mbps
		10	32.768 Mbps
		11	Reserved
2-1	RCVDATR		Transmit data rate
		00	8.192 Mbps
		01	16.384 Mbps
		10	32.768 Mbps
		11	Reserved

3.5 Clock and Frame Sync Selection and Polarity

Clock selection is independent for receive (RCVSRC) and transmit (XMTSRC) when CLKD is 1. RCVSRC and XMTSRC select both the source clock and frame sync for receive and transmit respectively, as shown in Table 3-4. All receive serial links use a common clock and frame sync. All transmit serial links use a common clock and frame sync.

Table 3-4   Clock and Frame Sync Source Selection (RCLKR and XCLKR)

Bit	Field	Value	Description
0	RCVSRC		Receive data clock and frame sync selection
		0	TCLK_A and TFS_A
		1	TCLK_B and TFS_B
0	XMTSRC		Transmit clock and frame sync selection
		0	TCLK_A and TFS_A
		1	TCLK_B and TFS_B

The polarity of the clock and frame sync is software selectable and can be defined independently for data and frame sync as well as receive and transmit. Table 3-5 provides a description of the clock polarity for data (RCVDCLKP and XMTDCLKP) and frame sync (RCVFCLKP and XMTFCLKP).

The polarity of the frame sync is software selectable and can be defined independently for receive and transmit. RCVFSYNCP and XMTFSYNCP determine the polarity of frame sync for receive and transmit respectively, as shown in Table 3-5.

Table 3-5   Frame Sync Polarity (RCR and XCR)

Bit	Field	Value	Description
7	RCVFSYNCP[1]		Receive frame sync polarity
		0	Receive frame sync is active low.
		1	Receive frame sync is active high.
7	XMTFSYNCP[2]		Transmit frame sync polarity
		0	Transmit frame sync is active low.
		1	Transmit frame sync is active high.

1. This field is write protected when SIU receive is enabled.
2. This field is write protected when SIU transmit is enabled.

3.6 Data Delay

The start of a frame (i.e., the frame boundary) is defined by the first clock cycle in which the frame synchronization signal is found to be active. The beginning of actual data reception or transmission with respect to the start of the frame can be delayed (or advanced) as required. This delay is called data delay.

Data delay is configured separately for receive and transmit data. The transmit data delay value stored in the register determines the delay, in sample clock periods, from the clock edge sampling the frame synchronization signal to the clock edge that starts driving the first data bit of the first timeslot. An XMTDATD value of 0 represents a delay of one clock period from the clock edge sampling the frame synchronization signal to the clock edge that starts driving the first data bit of the first timeslot when the same clock polarity is used for both events. If opposite clock edges are used for these two events, an XMTDATD value of 0 represents a delay of half a clock period from the clock edge sampling the frame synchronization signal to the clock edge that starts driving the first data bit. This calculation is shown in Table 3-6.

The receive data delay value stored in the register determines the delay, in sample clock periods, from the clock edge sampling the frame synchronization signal to the clock edge that samples the first data bit of the first timeslot. A RCVDATD value of 0 represents a delay of two clock periods from the clock edge sampling the frame synchronization signal to the clock edge that samples the first data bit of the first timeslot when the same clock polarity is used for both events. If opposite clock edges are used for these two events, a RCVDATD value of 0 represents a delay of one and a half clock periods from the clock edge sampling the frame synchronization signal to the clock edge that samples the first data bit. This calculation is also shown in Table 3-6.

Table 3-6   Data Delay (RCR and XCR)

Bit	Field	Value	Description
29-16	RCVDATD[1]	0x0000	Receive data delay [0-16383]
			n + 1.5 When clock polarity for data and frame sync differs
			n + 2 When clock polarity for data and frame sync is identical
29-16	XMTDATD[2]	0x0000	Transmit data delay [0-16383]
			n + 0.5 When clock polarity for data and frame sync differs
			n + 1 When clock polarity for data and frame sync is identical

1. This field is write protected when SIU receive is enabled.
2. This field is write protected when SIU transmit is enabled.

When compared with previous TDM interfaces in TI devices, there is a noticeable offset in the data delay value. This is the result of a full synchronous implementation. A data delay value of 0 in previous TDM implementations, like McBSP, required asynchronous logic that complicated the peripheral implementation. The paragraphs above stated the definition for data delay 0 for both receive and transmit links. Additional delay can be accommodated by increasing the data delay values. Smaller amounts of delay are also allowed. During 8.192 Mbps, single clock operation, there are exactly 1024 clock periods per frame. A data delay value of -1 is represented by a register value of 1023. A data delay value of -2 is represented by 1022. This same rationale is appropriate for 16.384 Mbps links and 32.768 Mbps links.

The following figures show examples of common TSIP timing implementations along with the associated data delay register settings. These figures can be used as templates when determining the proper settings in customer applications.

Figure 3-1   **Transmit Data Delay, FS Straddles Frame Boundary, Single-Rate**
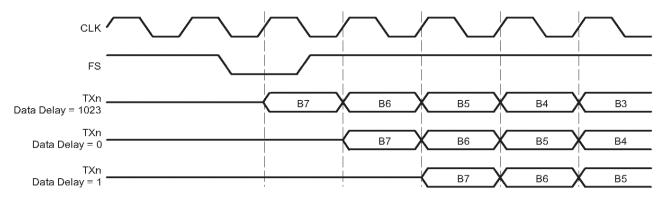
Figure 3-2   **Transmit Data Delay, FS Begins on Frame Boundary, Single-Rate**
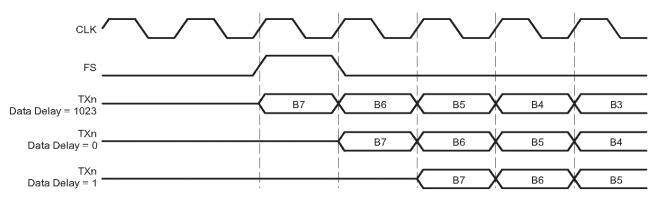
Figure 3-3   **Receive Data Delay, FS Straddles Frame Boundary, Single-Rate**
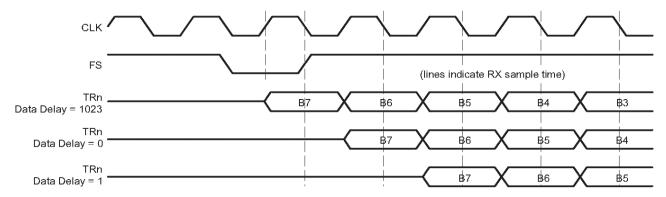
Figure 3-4   Receive Data Delay, FS Begins on Frame Boundary, Single-Rate
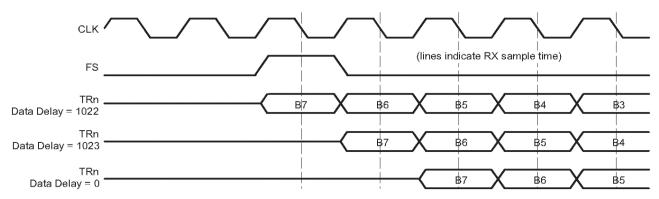
Figure 3-5   Transmit Data Delay, FS Straddles Frame Boundary, Double-Rate
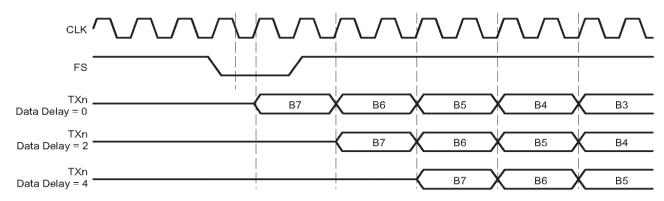
Figure 3-6   Transmit Data Delay, FS Begins on Frame Boundary, Double-Rate
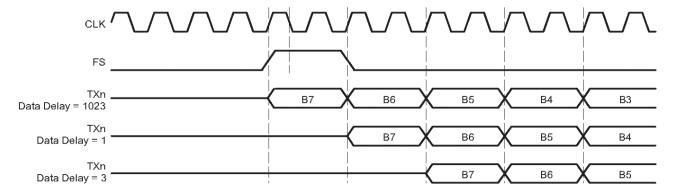
Figure 3-7   Receive Data Delay, FS Straddles Frame Boundary, Double-Rate
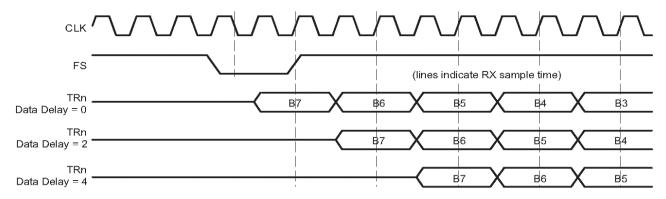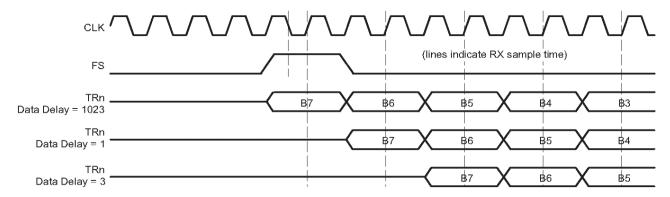
Figure 3-8   Receive Data Delay, FS Begins on Frame Boundary, Double-Rate

The TSIP interface supports both single-rate and double-rate clock modes. The offset calculations discussed above are appropriate for both single-rate and double-rate clock modes. The double-rate clock mode has much more flexibility to sample receive data streams and to offset transmit data streams. For 8.192 Mbps operation, there are 1024 clocks per frame in single-rate clock mode and 2048 clocks per frame in double-rate clock mode. Similarly, 16.384 Mbps operation has 2048 clocks per frame in single-rate clock mode and 4096 clocks per frame in double-rate clock mode and 32.768 Mbps operation has 4096 clocks per frame in single-rate clock mode and 8192 clocks per frame in double-rate clock mode. These values must be known when implementing a negative data delay in these different modes.

3.7 Data and Frames

3.7.1 Frame Size

A timeslot counter (10 bits) counts the timeslot events (or logical timeslots) that make up a frame. The RCVFSIZ bit in the Receive Size register (see Section 7.2.10) and the XMTFSIZ bit in the Transmit Size register (see Section 7.2.7) provide the rollover counts. The default rollover count for the timeslot counter is 127. In the default configuration the timeslot counter defines a frame size of 128 logical timeslots, which is the minimum number of timeslots for the frame. Because the 16.384 and 32.768 Mbps data rate options combine the shift registers in sets of two or four, there are two or four physical timeslots per logical timeslot. This typically means that the number of logical timeslots per frame is the same for all data rate options.

Table 3-7   Timeslot Counter (RSR and XSR)([1])

Bit	Field	Value	Description
6-0	RCVFSIZ([2])	0x7F	Receive frame size [0-127]
		n + 1	Number of logical timeslots per frame. Minimum number of logical timeslot periods per frame.
		≤6	Values of 6 and less are reserved. Attempts to write a value of 6 or less results in a value of 0x7F being registered.
6-0	XMTFSIZ([3])	0x7F	Transmit frame size [0-127]
		n + 1	Number of logical timeslots per frame. Minimum number of logical timeslot periods per frame.
		≤6	Values of 6 and less are reserved. Attempts to write a value of 6 or less results in a value of 0x7F being registered.

1. If CLKD is programmed to 0 or RCVSRC and XMTSRC will be written with the same value, then RCVFSIZ and XMTFSIZ must be programmed with the same value.
2. This field is write protected when SIU receive is enabled.
3. This field is write protected when SIU transmit is enabled.

3.7.2 Frame Count

A frame sync counter (8 bits) counts the frame sync input pulses (frame sync events) that make up a superframe. It outputs a superframe event on every Nth frame sync event. The RCVFCNT bit in the Receive register (see Section 7.2.10) and the XMTFCNT bit in the Transmit Size register (see Section 7.2.7) provide the rollover counts. The default rollover count for the frame sync counter is 3 which defines a superframe size of 4 frames. The frame sync counter is synchronized with the first valid frame sync input. The superframe event is provided to the TDMU, DMATCU, and interrupt generator.

Table 3-8   Frame Sync Counter (RSR and XSR)

Bit	Field	Value	Description
23-16	RCVFCNT([1])	0x3	Receive frame count [0-255]
		n + 1	Number of frames per superframe
23-16	XMTFCNT([2])	0x3	Transmit frame count [0-255]
		n + 1	Number of frames per superframe.

1. This field is write protected when SIU receive is enabled.
2. This field is write protected when SIU transmit is enabled.

3.8 Data Synchronization

3.8.1 Transmit and Receive Synchronization

The frame sync and time slot counter described in Section 3.7 is synchronized with the first valid frame sync input. Timeslot counter rollover event is used to switch from using the channel PING buffer to the channel PONG buffer for each channel or to switch from using the channel PONG buffer to the channel PING buffer for each channel. Channel buffer switching is locked together for all receive channels. It is also locked together for all transmit channels although the receive and transmit channels are not necessarily locked with each other.

For receive, this means that if the TDMU has been using the channel PING (or PONG) buffer while processing timeslots from frame N, it begins using the channel PONG (or PING) buffer while processing timeslots from frame N+1. If the DMATCU has not emptied the channel PONG buffer by the start of frame N+1, the TDMU does not overwrite the channel PONG buffer until it is emptied. Instead the data is discarded, frame N+1 is lost, and the buffer overrun flag is set.

For transmit, this means that if the transmit process logic has been using the channel PING (or PONG) buffer while processing timeslots for frame N it begins using the channel PONG (or PING) buffer while processing timeslots for frame N+1. If the DMATCU has not transferred any timeslot data into the channel PONG buffer by the start of frame N+1, there is no new data for the TDMU to process and send to the SIU. Due to the lack of data the TDMU sends an all ones (0xFF) value for all enabled timeslots. This means that frame N+1 went out with made-up data and the actual data for frame N+1 is never transmitted. The buffer underrun flag is set also.

3.8.2 Context Synchronization

Context switches between the A and B context definitions on each respective channel are made at the superframe completion. Each time a superframe event occurs the channel context for none, some, or all of the channels can be switched from A to B or from B to A. For the channels dedicated to the transmit serial links the superframe event occurs at the beginning of the Nth frame period. For the channels dedicated to the receive serial links the superframe event occurs at the end of the Nth frame period. In both transmit and receive the context switch takes place so that it is in effect for frame N+1.

Reset and Powerdown Mode

This chapter contains the details about:

4.1 SIU Reset and Shutdown Operation

The entire SIU can be reset with a single reset bit (SIURST), as shown in Table 4-1. Asserting the SIU reset causes both receive and transmit serial interfaces to halt operation.

Table 4-1   SIU Reset

Bit	Field	Value	Description
0	SIURST		SIU reset
			Reads:
		0	Reset is complete. SIU can be enabled for operation.
		1	Reset is in progress.
			Writes:
		0	Writes of 0 have no affect.
		1	Reset the SIU.

4.1.1 SIU Reset

When SIURST is set, the entire SIU module, transmit and receive, is reset. All registers that define SIU operation or provide the status for SIU operation are initialized to their predefined state. When read, SIURST indicates the status of the reset operation. SIURST remains 1 until the SIU has been shut down and all SIU control registers have been initialized to their predefined state. When the reset initialization is complete SIURST becomes 0. Writing 0 to SIURST has no effect.

4.1.2 Receive SIU Shutdown

When the receive interface is stopped, RCVENB is written as 0 after being 1, all of the serial-to parallel registers, data registers, and counters used for the receive operation are cleared. The programmed configuration defining the operation of the receive interface is not affected when the receive interface is disabled. No events or interrupts are generated while the receive interface is disabled. When RCVENB is written as 1 after being 0 the operation of the receive interface is enabled, the first active frame sync input is recognized and serial data capture begins immediately. The first active frame sync is recognized only when the enabled SIU receive is able to first detect the frame sync input in an inactive state.

4.1.3 Transmit SIU Shutdown

When the transmit interface is stopped, XMTENB is written as 0 after being 1, all of the parallel to-serial registers, data registers, and counters used for the transmit operation are cleared. The output state of the output buffers is controlled by the output state configuration for disabled timeslots. The programmed configuration defining the operation of the transmit interface is not affected when the transmit interface is disabled. No events or interrupts are generated while the transmit interface is disabled. When XMTENB is written as 1 after being 0 the operation of the transmit interface is enabled, the first active frame sync input is recognized but data transmission does not begin until the second active frame sync input is recognized. This assures sufficient time for the serial interface to signal the DMA controller and for the DMA controller to respond with timeslot 0 data prior to the second active frame sync pulse. The first active frame sync is recognized only when the enabled SIU transmit is able to first detect the frame sync input in an inactive state. During the period between the first and second active frame syncs the data output is in the disabled state, as defined in Section 2.3, for all timeslots.

4.2 TDMU and DMATCU Software Reset and Shutdown Operation

The TDMU and DMATCU can be reset or stopped (disabled) entirely. Individual channels can be stopped (disabled) independently as well. Resetting the entire TDMU and DMATCU causes all channels to be placed into a stopped (disabled) state. Stopping the TDMU and DMATCU causes the operation to be stopped but does not necessarily alter the state of the individual channels.

Stopping individual channels results in those channels not responding to the timeslot, frame, and superframe events generated by the serial interface. Stopping the TDMU channels does not affect the operation of the SIU.

Reset of the TDMU and DMATCU occurs when DMARST is written as a 1. There is one DMAENB bit for the TDMU and DMATCU together. There are individual XMTCHnENB bits and RCVCHnENB bits, one for each channel.

4.3 Powerdown

The TSIP supports clock stop handshaking as specified in CBA3.0 and implemented with the use of the local powerdown and sleep controller (LPSC).

The LPSC can be configured to make clock stop requests to the TSIP. If the TSIP is in an idle state the clock stop request is acknowledged and the gated clock can then be stopped. If TSIP is in an active state the clock stop request is not acknowledged.

Emulation and Test Modes

This chapter contains the details about:

5.1 Emulation and Test Modes

Two control bits, FREE and SOFT, determine the response of the serial interface to the emulation input. Table 5-1 provides the definitions of FREE and SOFT.

Table 5-1   Emulation Control (ETR)

Bit	Field	Value	Description
0	FREE(1)		Free run selection
		0	SOFT takes effect.
		1	Ignore emulation input. Continue to function normally.
1	SOFT(1)		Halt selection
		0	Halt receive and transmit SIU, TDMU, and DMATCU at the completion of the current frame.
		1	Halt receive TDMU and DMATCU at the completion of the current frame. Transmit continues to function normally.

1. This field is write protected once DMA is enabled.

5.1.1 Emulation -Free Run

When FREE=1 it defines the operation to be free running even under emulation halt conditions. The TSIP continues to operate as normal with data being received, buffered, and transferred to memory and data being transferred from memory, buffered, and transmitted based on the presence of the necessary external serial data clock and frame sync inputs and the internal clock input.

5.1.2 Emulation -Halt

When FREE=0 the operation is defined by SOFT. The presence of the necessary external serial data clock and frame sync inputs and the internal clock input is assumed as well.

For SOFT=0, the operation of the TSIP continues as normal until the end of the current frame. At the end of the current frame the TSIP operation gets halted. The frame interrupt/event for the beginning of the next frame and the possible superframe interrupt/event is generated.

For SOFT=1, the operation of the TSIP continues as normal until the end of the current frame. Receive operation is halted from very next frame, transmit continues to work normally in this mode. But as the DMA controller is not functional, transmit memory buffers are not updated during this time, since the same data for the sequence of frames that is currently in the memory buffer continues to be sent over and over.

5.1.3 Test Modes

Internal test modes are provided to test the functionality of the TSIP. These test modes allow the functionality to be tested with a minimum of external stimulus or with a minimal amount of internal software to support data processing.

5.1.3.1 Serial Interface Test Modes

The TSIP SIU offers operational test supports data loop back (DLB) and link loop back (LLB), as illustrated in Figure 5-1. The test modes, when enabled, are mutually exclusive. Control bits, SIUTST, and LBS determine the operation of the serial interface for test. The definitions of SIUTST and LBS are found in Table 5-2. A common external clock and frame sync for receive and transmit are required for loop back testing to be successful.

Figure 5-1   Loopback Test Modes
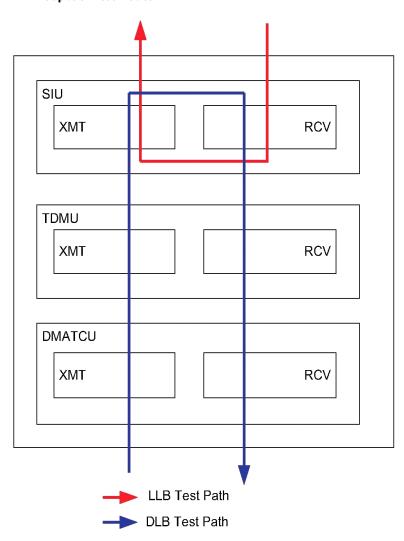
Table 5-2   SIU Test Control (ETR)

Bit	Field	Value	Description
16	SIUTST[1]		Serial test mode enable
		0	Normal operating mode is enabled. LBS is ignored.
		1	Serial loopback test is enabled.
17	LBS[1]		Serial test mode selection
		0	DLB is selected.
		1	LLB is selected.

1. This field is write protected once DMA is enabled.

5.1.3.2 DLB Testing

DLB enables the serial link interface to be tested from inside the device. In this test configuration the SXn outputs are internally connected to the corresponding SRn inputs. The TXn output pins are disabled. The TRn input pins are ignored. Corresponding timeslots for transmit and receive should be enabled in the TSIP DMA. Successful operation is determined by comparing a receive frame buffer in memory with the corresponding transmit frame buffer in memory.

5.1.3.3 LLB Testing

LLB enables the serial link interface to be tested from outside the device. In this test configuration the DRn outputs are internally connected to the corresponding DXn inputs. Successful operation is determined by comparing the transmit serial links with a delayed version of the receive serial links.

TSIP Initialization Procedure
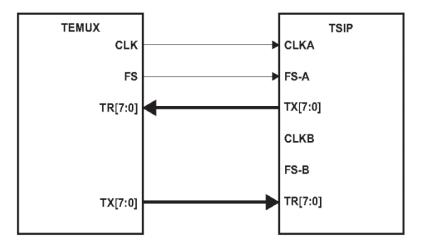
This chapter discusses initialization procedure.

6.1 TSIP Initialization Procedure

Configure the TSIP data rate, clock settings, frame sync settings, and the DMA channel configurations based on the system configuration. Figure 6-1 is an example of a typical case where the TSIP is used in a system configuration in which it is connected to a TEMUX device with 8-Mbps data rate and 16-MHz clock frequency, with the same clock and frame sync being used for transmit and receive.

Figure 6-1   TSIP Connected to TEMUX

In this configuration, CLKA and FSA are used by the transmit and receive logic for clock and frame synchronization.

1. Ensure that XMTENB, RCVENB, DMA channel enable, and DMAENB are cleared to make certain that the TSIP is in reset. Reset the SIU and the DMA to ensure that the TSIP is properly shut down.

2. As the same clock and frame sync is being used by the TSIP TX and RX, configure CLKD to select the redundant clock.

3. If redundant clock mode is used, XMTSRC, XMTCLKM and XMTDATR need to be configured to select the clock source and clock rate. If redundant clock mode is not used, configure both transmit and receive registers for the same. The XMTSRC field has to be cleared for the TSIP to use the clock source A and FSA. For 8 Mbps data rate, XMTDATR has to be set to 0x0. As the TSIP data clock rate is at 16 MHz for this system, XMTCLKM sound be configured to 0x0 for selecting the double data rate clock mode.

4. Clock and frame sync polarities have to be configured depending on whether the TEMUX device is driving active low or active high signal for frame sync and depending on the clock-edge at which it is driving the data.

5. Configure the XMTDIS state depending on how the transmit lanes have to be tri-stated in the system.

6. XMTDATD and RCVDATD needs to be set based on the delay in the system for the first valid data to be received by the TSIP after the frame sync is detected. SIU is set to transmit/receive the data on its TX and TR lanes.

7. Based on the call setup and teardown messages from the host, the system should be configured for the transmit and receive frame size after determining the number of logical timeslots that need to be enabled in a frame. Frame count in each super frame depends on the application.

8.   Configure the endianness and priority of the TSIP master in the system as per the system requirements.

9.   Enable the interrupts required to the system from the TSIP.

10. Configure the DMA channel context registers depending on the channels to be enabled. Once DMA context is written, configure the CnID and enable the channel by writing into the CHnENB bit.

11. Enable the DMA by setting DMAENB and enable SIU by setting the XMTENB and RCVENB bits.

12. The TSIP is set to receive and transmit data on the next Frame Sync.

13. The TSIP switches between the PING and PONG buffers on every superframe. Hence, the system can configure the DMA channel parameters for a context switch while the TSIP transfer is going on. Bitmap slots can be re-enabled depending on the call setup messages received for the next superframe using the ping-pong buffer mechanism.

Registers

This section includes the TSIP memory map information and registers.

7.1 TSIP Memory Map

Table 7-1 to Table 7-13 provide additional register memory map information.
Table 1-1 lists the summary memory map for the registers for the TSIP.

Table 7-1   TSIP Module Registers

Offset Address	Acronym	Type	Reset Value	Reset By (¹)	Description
0x00000000	PID	Constant	0x600101	H	PID Register
0x00000004	EMUTST	Configuration	0x00000001	H, S	Emulation and Test Register
0x00000008	RST	Configuration	0x00000000	H	Reset Register
0x0000000C-0x0000007C	-	-	-	-	Reserved

1. H=Hardware reset, S=SIURST, D=DMARST

Table 7-2   Serial Interface Registers

Offset Address	Acronym	Type	Reset Value	Reset By (¹)	Description
0x00000080	SIUCTL	Configuration	0x00000000	H, S	SIU Global Control Register
0x00000084-0x0000009C	-	-	-	-	Reserved
0x000000A0	XCLK	Configuration	0x00000000	H, S	Transmit Clock Source Register
0x000000A4	XCTL	Configuration	0x00000000	H, S	Transmit Control Register
0x000000A8	XSIZE	Configuration	0x0000007F	H, S	Transmit Size Register
0x0000000AC-0x000000BC	-	-	-	-	Reserved
0x0000000C0	RCLK	Configuration	0x00000000	H, S	Receive Clock Source Register
0x000000C4	RCTL	Configuration	0x00000000	H, S	Receive Control Register
0x000000C8	RSIZE	Configuration	0x0000007F	H, S	Receive Size Register
0x0000000CC-0x000000FC	-	-	-	-	Reserved

1. H=Hardware reset, S=SIURST, D=DMARST

Table 7-3   TDMU Global Registers

Offset Address	Acronym	Type	Reset Value	Reset By (¹)	Description
0x00000100	TDMUCTL	Configuration	0x00000000	H, D	TDMU Global Control Register
0x00000104	XFRFC	Status	0x00000000	H, D	Transmit Free Running Frame Counter
0x00000108	RFRFC	Status	0x00000000	H, D	Receive Free Running Frame Counter
0x0000010C	TDMUCFG	Configuration	0x00000000	H, D	TDMU Global Configuration Register
0x00000110	XBMST	Status	0x00000000	H, D	Transmit Channel Bitmap Active Status Register
0x00000114	RBMST	Status	0x00000000	H, D	Receive Channel Active Status Register
0x00000118 - 0x0000017C	-	-	-	-	Reserved

1. H=Hardware reset, S=SIURST, D=DMARST

Table 7-4   DMATCU Global Registers

Offset Address	Acronym	Type	Reset Value	Reset By (1)	Description
0x00000180	DMACTL	Configuration	0x00000077	H, D	DMATCU Global Control Register
0x00000184	XDLY	Configuration	0x0000007F	H, D	Transmit Timeslot Delay Counter
0x00000188	RDLY	Configuration	0x0000007F	H, D	Receive Timeslot Delay Counter
0x0000018C	-	-	-	-	Reserved
0x00000190	XCHST	Status	0x00000000	H, D	Transmit Channel Configuration Active Status Register
0x00000194	RCHST	Status	0x00000000	H, D	Receive Channel Active Status Register
0x00000198 - 0x000001FC	-	-	-	-	Reserved

1. H=Hardware reset, S=SIURST, D=DMARST

Table 7-5   TX/RX Channel Error Registers

Offset Address	Register Description
0x00000200 - 0x0000020C	TX/RX channel 0 error log registers
0x00000210 - 0x0000021C	TX/RX channel 1error log registers
0x00000220 - 0x0000022C	TX/RX channel 2 error log registers
0x00000230 - 0x0000023C	TX/RX channel 3 error log registers
0x00000240 - 0x0000024C	TX/RX channel 4 error log registers
0x00000250 - 0x0000025C	TX/RX channel 5 error log registers
0x00000260 - 0x000003FC	Reserved

Table 7-6   TDMU Channel 0 Error Registers

Offset Address	Acronym	Type	Reset Value	Reset By(1)	Description
0x00000200	ERRCTL0	Configuration	0x00000000	H, D	TX/RX Channel 0 Error Control Register
0x00000204	ERRCNT0	Status	0x00000000	H, D	TX/RX Channel 0 Error Count Register
0x00000208	ERRQ0	Status	0x00000000	H, D	TX/RX Channel 0 Error Queue Register
0x0000020C	-	-	-	-	Reserved

1. H=Hardware reset, S=SIURST, D=DMARST

Table 7-7   TDMU Channel Registers (Part 1 of 2)

Offset Address	Register Description
0x00000800 - 0x0000081C	Transmit channel 0 registers
0x00000820 - 0x0000083C	Transmit channel 1 registers
0x00000840 - 0x0000085C	Transmit channel 2 registers
0x00000860 - 0x0000087C	Transmit channel 3 registers
0x00000880 - 0x0000089C	Transmit channel 4 registers
0x000008A0 - 0x000008BC	Transmit channel 5 registers
0x000008C0 - 0x00000BFC	Reserved
0x00000C00 - 0x00000C1C	Receive channel 0 registers
0x00000C20 - 0x00000C3C	Receive channel 1 registers
0x00000C40 - 0x00000C5C	Receive channel 2 registers

Table 7-7   TDMU Channel Registers (Part 2 of 2)

Offset Address	Register Description
0x00000C60 - 0x00000C7C	Receive channel 3 registers
0x00000C80 - 0x00000C9C	Receive channel 4 registers
0x00000CA0 - 0x00000CBC	Receive channel 5 registers
0x00000CC0 - 0x00000FFC	Reserved

Table 7-8   TDMU Transmit Channel 0 Registers

Offset Address[1]	Acronym	Type	Reset Value	Reset By [2]	Description
0x00000800	XCHEN0	Configuration	0x00000000	H, D	Transmit Channel 0 Enable Register
0x00000804 - 0x0000081C	-	-	-	-	Reserved

1. Registers for Transmit Channels 1-5 are equivalent with offsets as defined in Table 31.

2. H=Hardware reset, S=SIURST, D=DMARST

Table 7-9   TDMU Receive Channel 0 Registers

Offset Address[1]	Acronym	Type	Reset Value	Reset By [2]	Description
0x00000C00	RCHEN0	Configuration	0x00000000	H, D	Receive Channel 0 Enable Register
0x00000C04 - 0x00000C1C	-	-	-	-	Reserved

1. Registers for Receive Channels 1-5 are equivalent with offsets as defined in Table 31.

2. H=Hardware reset, S=SIURST, D=DMARST

Table 7-10   DMATCU Channel Registers (Part 1 of 2)

Offset Address	Register Description
0x00001000 - 0x0000101C	Transmit channel 0 configuration A registers
0x00001020 - 0x0000103C	Transmit channel 0 configuration B registers
0x00001040 - 0x0000105C	Transmit channel 1 configuration A registers
0x00001060 - 0x0000107C	Transmit channel 1 configuration B registers
0x00001080 - 0x0000109C	Transmit channel 2 configuration A registers
0x000010A0 - 0x000010BC	Transmit channel 2 configuration B registers
0x000010C0 - 0x000010DC	Transmit channel 3 configuration A registers
0x000010E0 - 0x000010FC	Transmit channel 3 configuration B registers
0x00001100 - 0x0000111C	Transmit channel 4 configuration A registers
0x00001120 - 0x0000113C	Transmit channel 4 configuration B registers
0x00001140 - 0x0000115C	Transmit channel 5 configuration A registers
0x00001160 - 0x0000117C	Transmit channel 5 configuration B registers
0x00001180 - 0x000017FC	Reserved
0x00001800 - 0x0000181C	Receive channel 0 configuration A registers
0x00001820 - 0x0000183C	Receive channel 0 configuration B registers
0x00001840 - 0x0000185C	Receive channel 1 configuration A registers
0x00001860 - 0x0000187C	Receive channel 1 configuration B registers
0x00001880 - 0x0000189C	Receive channel 2 configuration A registers
0x000018A0 - 0x000018BC	Receive channel 2 configuration B registers
0x000018C0 - 0x000018DC	Receive channel 3 configuration A registers
0x000018E0 - 0x000018FC	Receive channel 3 configuration B registers
0x00001900 - 0x0000191C	Receive channel 4 configuration A registers

Table 7-10   DMATCU Channel Registers (Part 2 of 2)

Offset Address	Register Description
0x00001920 - 0x0000193C	Receive channel 4 configuration B registers
0x00001940 - 0x0000195C	Receive channel 5 configuration A registers
0x00001960 - 0x0000197C	Receive channel 5 configuration B registers
0x00001980 - 0x00001FFC	Reserved
End of Table 7-10	

Table 7-11   DMATCU Transmit Channel 0 Registers

Offset Address[1]	Acronym	Type	Reset Value	Reset By ([2])	Description
0x00001000	DXCH_ABASE0	Configuration	0x00000000	H, D	Transmit Channel 0 Memory Base Address Register A
0x00001004	DXCH_AFALLOC0	Configuration	0x00000000	H, D	Transmit Channel 0 Frame Allocation Register A
0x00001008	DXCH_AFSIZE0	Configuration	0x00000000	H, D	Transmit Channel 0 Frame Size Register A
0x0000100C	DXCH_AFCNT0	Configuration	0x00000000	H, D	Transmit Channel 0 Frame Count Register A
0x000001010 - 0x0000101C	-	-	-	-	Reserved
0x000001020	DXCH_BBASE0	Configuration	0x00000000	H, D	Transmit Channel 0 Memory Base Address Register B
0x000001024	DXCH_BFALLOC0	Configuration	0x00000000	H, D	Transmit Channel 0 Frame Allocation Register B
0x000001028	DXCH_BFSIZE0	Configuration	0x00000000	H, D	Transmit Channel 0 Frame Size Register B
0x00000102C	DXCH_BFCNT0	Configuration	0x00000000	H, D	Transmit Channel 0 Frame Count Register B
0x000001030 - 0x00000103C	-	-	-	-	Reserved
End of Table 7-11					

1. Registers for Transmit channels 1-5 are equivalent with offsets as defined in .

2. H=Hardware reset, S=SIURST, D=DMARST

Table 7-12   DMATCU Receive Channel 0 Registers (Part 1 of 2)

Offset Address[1]	Acronym	Type	Reset Value	Reset By ([2])	Description
0x00001800	DRCH_ABASE0	Configuration	0x00000000	H, D	Receive Channel 0 Memory Base Address Register A
0x00001804	DRCH_AFALLOC0	Configuration	0x00000000	H, D	Receive Channel 0 Frame Allocation Register A
0x00001808	DRCH_AFSIZE0	Configuration	0x00000000	H, D	Receive Channel 0 Frame Size Register A
0x0000180C	DRCH_AFCNT0	Configuration	0x00000000	H, D	Receive Channel 0 Frame Count Register A
0x000001810 - 0x0000181C	-	-	-	-	Reserved
0x000001820	DRCH_BBASE0	Configuration	0x00000000	H, D	Receive Channel 0 Memory Base Address Register B
0x000001824	DRCH_BFALLOC0	Configuration	0x00000000	H, D	Receive Channel 0 Frame Allocation Register B
0x000001828	DRCH_BFSIZE0	Configuration	0x00000000	H, D	Receive Channel 0 Frame Size Register B

Table 7-12   DMATCU Receive Channel 0 Registers (Part 2 of 2)

Offset Address[1]	Acronym	Type	Reset Value	Reset By (²)	Description
0x00000182C	DRCH_BFCNT0	Configuration	0x00000000	H, D	Receive Channel 0 Frame Count Register B
0x000001830 - 0x0000183C	-	-	-	-	Reserved
End of Table 7-12					

1. Registers for Receive channels 1-5 are equivalent with offsets as defined in Table 7-10.

2. H=Hardware reset, S=SIURST, D=DMARST

Table 7-13   TDMU Channel Bitmap Registers

Offset Address	Acronym	Register Description
0x00008000 - 0x000080FC	XBM_XBMA0	Transmit channel 0 bitmap A
0x00008100 - 0x000081FC	XBM_XBMB0	Transmit channel 0 bitmap B
0x00008200 - 0x000082FC	XBM_XBMA1	Transmit channel 1 bitmap A
0x00008300 - 0x000083FC	XBM_XBMB1	Transmit channel 1 bitmap B
0x00008400 - 0x000084FC	XBM_XBMA2	Transmit channel 2 bitmap A
0x00008500 - 0x000085FC	XBM_XBMB2	Transmit channel 2 bitmap B
0x00008600 - 0x000086FC	XBM_XBMA3	Transmit channel 3 bitmap A
0x00008700 - 0x000087FC	XBM_XBMB3	Transmit channel 3 bitmap B
0x00008800 - 0x000088FC	XBM_XBMA4	Transmit channel 4 bitmap A
0x00008900 - 0x000089FC	XBM_XBMB4	Transmit channel 4 bitmap B
0x00008A00 - 0x00008AFC	XBM_XBMA5	Transmit channel 5 bitmap A
0x00008B00 - 0x00008BFC	XBM_XBMB5	Transmit channel 5 bitmap B
0x00008C00 - 0x0000BFFC	-	Reserved
0x0000C000 - 0x0000C0FC	RBM_RBMA0	Receive channel 0 bitmap A
0x0000C100 - 0x0000C1FC	RBM_RBMB0	Receive channel 0 bitmap B
0x0000C200 - 0x0000C2FC	RBM_RBMA1	Receive channel 1 bitmap A
0x0000C300 - 0x0000C3FC	RBM_RBMB1	Receive channel 1 bitmap B
0x0000C400 - 0x0000C4FC	RBM_RBMA2	Receive channel 2 bitmap A
0x0000C500 - 0x0000C5FC	RBM_RBMB2	Receive channel 2 bitmap B
0x0000C600 - 0x0000C6FC	RBM_RBMA3	Receive channel 3 bitmap A
0x0000C700 - 0x0000C7FC	RBM_RBMB3	Receive channel 3 bitmap B
0x0000C800 - 0x0000C8FC	RBM_RBMA4	Receive channel 4 bitmap A
0x0000C900 - 0x0000C9FC	RBM_RBMB4	Receive channel 4 bitmap B
0x0000CA00 - 0x0000CAFC	RBM_RBMA5	Receive channel 5 bitmap A
0x0000CB00 - 0x0000CBFC	RBM_RBMB5	Receive channel 5 bitmap B
0x0000CC00 - 0x0000FFFC	-	Reserved
End of Table 7-13		

Table 7-14   TDMU Channel Buffers (Part 1 of 3)

Offset Address	Acronym	Register Description
0x00010000 - 0x000100FC	XPING0	Transmit Channel 0 PING Buffer
0x00010100 - 0x000003FC	-	Reserved
0x00010400 - 0x000004FC	XPONG0	Transmit Channel 0 PONG Buffer

Table 7-14   TDMU Channel Buffers (Part 2 of 3)

Offset Address	Acronym	Register Description
0x00010500 - 0x000007FC	-	Reserved
0x00010800 - 0x000008FC	XPING1	Transmit Channel 1 PING Buffer
0x00010900 - 0x00000BFC	-	Reserved
0x00010C00 - 0x00000CFC	XPONG1	Transmit Channel 1 PONG Buffer
0x00010D00 - 0x00000FFC	-	Reserved
0x00011000 - 0x00001DFC	XPING2	Transmit Channel 2 PING Buffer
0x00011100 - 0x000013FC	-	Reserved
0x00011400 - 0x000014FC	XPONG2	Transmit Channel 2 PONG Buffer
0x00011500 - 0x000017FC	-	Reserved
0x00011800 - 0x000018FC	XPING3	Transmit Channel 3 PING Buffer
0x00011900 - 0x00001BFC	-	Reserved
0x00011C00 - 0x00001CFC	XPONG3	Transmit Channel 3 PONG Buffer
0x00011D00 - 0x00001FFC	-	Reserved
0x00012000 - 0x00002DFC	XPING4	Transmit Channel 4 PING Buffer
0x00012100 - 0x000023FC	-	Reserved
0x00012400 - 0x000024FC	XPONG4	Transmit Channel 4 PONG Buffer
0x00012500 - 0x000027FC	-	Reserved
0x00012800 - 0x000028FC	XPING5	Transmit Channel 5 PING Buffer
0x00012900 - 0x00002BFC	-	Reserved
0x00012C00 - 0x00002CFC	XPONG5	Transmit Channel 5 PONG Buffer
0x00012D00 - 0x00007FFC	-	Reserved
0x00018000 - 0x00008DFC	RPING0	Receive Channel 0 PING Buffer
0x00018100 - 0x000083FC	-	Reserved
0x00018400 - 0x000084FC	RPONG0	Receive Channel 0 PONG Buffer
0x00018500 - 0x000087FC	-	Reserved
0x00018800 - 0x000088FC	RPING1	Receive Channel 1 PING Buffer
0x00018900 - 0x00008BFC	-	Reserved
0x00018C00 - 0x00008CFC	RPONG1	Receive Channel 1 PONG Buffer
0x00018D00 - 0x00008FFC	-	Reserved
0x00019000 - 0x00009DFC	RPING2	Receive Channel 2 PING Buffer
0x00019100 - 0x000093FC	-	Reserved
0x00019400 - 0x000094FC	RPONG2	Receive Channel 2 PONG Buffer
0x00019500 - 0x000097FC	-	Reserved
0x00019800 - 0x000098FC	RPING3	Receive Channel 3 PING Buffer
0x00019900 - 0x00009BFC	-	Reserved
0x00019C00 - 0x00009CFC	RPONG3	Receive Channel 3 PONG Buffer
0x00019D00 - 0x00009FFC	-	Reserved
0x0001A000 - 0x0000ADFC	RPING4	Receive Channel 4 PING Buffer
0x0001A100 - 0x0000A3FC	-	Reserved
0x0001A400 - 0x0000A4FC	RPONG4	Receive Channel 4 PONG Buffer
0x0001A500 - 0x0000A7FC	-	Reserved
0x0001A800 - 0x0000A8FC	RPING5	Receive Channel 5 PING Buffer
0x0001A900 - 0x0000ABFC	-	Reserved

Table 7-14   TDMU Channel Buffers (Part 3 of 3)

Offset Address	Acronym	Register Description
0x0001AC00 - 0x0000ACFC	RPONG5	Receive Channel 5 PONG Buffer
0x0001AD00 - 0x0000FFFC	-	Reserved
End of Table 7-14		

7.2 Register Description

This section provides a detailed definition of the peripheral module constant registers.

7.2.1 PID Register

The PID register is shown in Figure 7-1 and described in Table 7-15.

Figure 7-1   PID Register (0x00000000)

31	16	15	8	7	0
MODID		Major		MINOR	
R-0x0060		R-0x01		R-1x00	

LEGEND: R/W = Read/Write; R = Read only; -n = Value after reset

Table 7-15   PID Register Field Descriptions

Bit	Field	Value	Description
31-16	MODID	0x0060	Module ID
15-8	MAJOR	0x01	Major Version Value
7-0	MINOR	0x00	Minor Version Value

7.2.2 Emulation and Test Register

The emulation and test register is shown in Figure 7-2 and described in Table 7-16.

Figure 7-2   PID Register (0x00000000)

31	18	17	16
Reserved		LBS	SIUTST
R-0x00		R/W-0	R/W-0

15	2	1	0
Reserved		SOFT	FREE
R-0x00		R/W-0	R/W-1

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-16   Emulation and Test Register Field Descriptions (Part 1 of 2)

Bit	Field	Value	Description
31-18	Reserved	0x0000	Reserved
17	LBS[1]		Serial test mode selection
		0	Internal
		1	External
16	SIUTST[1]		Serial test mode enable
		0	Normal operating mode is enabled. LBS is ignored.
		1	Serial loopback test mode is enabled.

Table 7-16   Emulation and Test Register Field Descriptions (Part 2 of 2)

Bit	Field	Value	Description
15-2	Reserved	0x00	Reserved
1	SOFT([1])		Halt selection
		0	Halt SIU, TDMU, and DMTCU at the completion of the current frame.
		1	Halt receive TDMU and DMATCU at the completion of the current frame.
0	Free([1])		Free run selection
		0	SOFT takes effect.
		1	Ignore emulation input. Continue to function normally.

1. This field is write protected once DMA is enabled.

7.2.3 Reset Register

The reset register is shown in Figure 7-3 and described in Table 7-17.

Figure 7-3   Reset Register (0x00000008)

31		2	1	0
Reserved			DMARST	SIURST
R-0x00000000			R/W-0	R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-17   Reset Register Field Descriptions

Bit	Field	Value	Description
31-2	Reserved	0x00000000	Reserved
1	DMARST		TDMU and DMATCU reset
			Reads:
		0	Reset is complete. TDMU and DMATCU can be enabled for operation.
		1	Reset is in progress.
			Writes:
		0	Writes of 0 have no affect.
		1	Reset the TDMU and DMATCU.
0	SIURST		SIU reset
			Reads:
		0	Reset is complete. SIU can be enabled for operation.
		1	Reset is in progress.
			Writes:
		0	Writes of 0 have no affect.
		1	Reset the SIU.

7.2.4 SIU Global Control Register

The SIU global control register is shown in Figure 7-4 and described in Table 7-18.

Figure 7-4   SIU Global Control Register (0x00000080) (Part 1 of 2)

31	16
Reserved	
R-0x00000000	

Figure 7-4   SIU Global Control Register (0x00000080) (Part 2 of 2)

15					5	4	3	2	1	0
Reserved						CLKD	Reserved		RCVENB	XMTENB
R-0x00000000						R/W-0	R-00		R/W-0	R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset
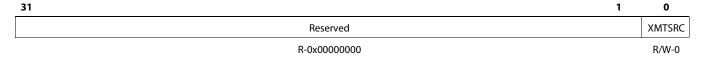
Table 7-18   SIU Global Control Register Field Descriptions

Bit	Field	Value	Description
31-5	Reserved	0x00000000	Reserved
4	CLKD		Clock redundancy selection
		0	Redundant clock mode. Transmit and receive share the same clock and frame sync. The common selections are those defined by SMTSRC, SMTCLKM, and SMTDATR in the transmit select and enable register. The RCVSRC, RCVCLKM, and RCVDATR fields are ignored.
		1	Dual clock mode. Transmit and receive are specified totally independently.
3-2	Reserved	0	Reserved
1	RCVENB		Receive enable
		0	Receive serial interface is disabled/stopped.
		1	Receive serial interface is enabled/running.
0	XMTENB		Transmit enable
		0	Transmit serial interface is disabled/stopped.
		1	Transmit serial interface is enabled/running.

7.2.5 Transmit Clock Source Register

The transmit clock source register is shown in Figure 7-5 and described in Table 7-19.

Figure 7-5   Transmit Clock Source Register (0x000000A0)

31	1	0
Reserved		XMTSRC
R-0x00000000		R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-19   Transmit Clock Source Register Field Descriptions

Bit	Field	Value	Description
31-1	Reserved	0x00000000	Reserved
0	XMTSRC		Transmit clock and frame sync selection
		0	TCLK_A and TFS_A
		1	TCLK_B and TFS_B

7.2.6 Transmit Control Register

The transmit control register is shown in Figure 7-6 and described in Table 7-20.

Figure 7-6   Transmit Control Register (0x000000A4)

31	30	29												16
Reserved		XMTDATD												

R-0x0	R/W-0x0000

15										11	10	9	8
Reserved											XMTDLY	XMTDIS	

R-0x00	R/W-0	R/W-00

7	6	5	4	3	2	1	0
XMTFSYNCP	XMTFCLKP	XMTDCLKP	Reserved		XMTDATR		XMTCLKM

R/W-0	R/W-0	R/W-0	R-00	R/W-00	R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset
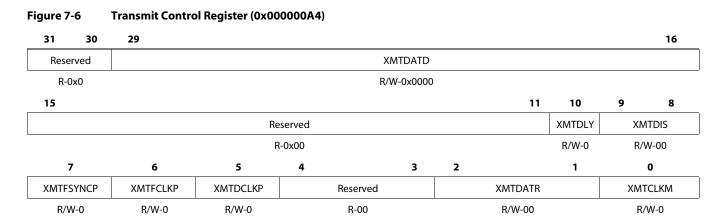
Table 7-20   Transmit Control Register Field Descriptions(1) (Part 1 of 2)

Bit	Field	Value	Description
31-30	Reserved	0x0	Reserved
29-16	XMTDATD(2)	0x0000	Transmit data delay [0-16383]
		n + 0.5	When clock polarity for data and frame sync differs.
		n + 1	When clock polarity for data and frame sync is identical.
15-11	Reserved	0x00	Reserved
10	XMTDLY(2)		Transmit output enable delay. Defines fixed delay for enabling the output buffer on the first bit of an enabled timeslot that immediately follows a disabled timeslot. This control applies only when the disable state for the output buffer is high impedance.
		0	No added delay.
		1	Output enable is delayed by one-half serial clock period.
9-8	XMTDIS(2)		Transmit output disable state
		00	High impedance
		01	Reserved
		10	Driven low
		11	Driven high
7	XMTFSYNCP(2)		Transmit frame sync polarity
		0	Transmit frame sync is active low.
		1	Transmit frame sync is active high.
6	XMTFCLKP(2)		Transmit frame sync clock polarity
		0	Transmit frame sync is sampled on the rising edge of TCLK_x.
		1	Transmit frame sync is sampled on the falling edge of TCLK_x.
5	XMTDCLKP(2)		Transmit data clock polarity
		0	Transmit data is driven on the rising edge of TCLK_x.
		1	Transmit data is driven on the falling edge of TCLK_x.
4-3	Reserved	0	Reserved
2-1	XMTDATR(2)		Transmit data rate
		00	8.192 Mbps
		01	16.384 Mbps
		10	32.768 Mbps
		11	Reserved

Table 7-20   Transmit Control Register Field Descriptions(¹) (Part 2 of 2)

Bit	Field	Value	Description
0	XMTCLKM(²)		Transmit clock mode
		0	TCLK_x is a double rate clock.
		1	TCLK_x is a single rate clock.

1. If CLKD is programmed to 0 or RCVSRC and XMTSRC will be written with the same value, then RCVFSYNCP and XMTFSYNCP must be programmed with the same value. Likewise, RCVFCLKP and XMTFCLKP must be programmed with the same value.
2. This field is write protected when SIU transmit is enabled.

7.2.7 Transmit Size Register

The transmit size register is shown in Figure 7-7 and described in Table 7-21.

Figure 7-7   Transmit Size Register (0x000000A8)

31	24	23	16	15	7	6	0
Reserved		XMTFCNT		Reserved		XMTFSIZ	
R-0x000		R/W-0x3		R-0x00		R/W-0x07F	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-21   Transmit Size Register Field Descriptions(¹)

Bit	Field	Value	Description
31-24	Reserved	0x000	Reserved
23-16	XMTFCNT(²)	0x3	Transmit frame count [0-255]
			$n + 1$ Number of frames per superframe
15-7	Reserved	0x00	Reserved
6-0	XMTFSIZ(²)	0x7F	Transmit frame size [0-27]
			$n + 1$ Number of logical timeslots per frame. Minimum number of logical timeslot periods per frame.
			≤6 Values of 6 and less are reserved. Attempts to write a value of 6 or less result in a value of 0x7F being registered.

1. If CLKD is programmed to 0 or RCVSRC and XMTSRC will be written with the same value, then RCVFSIZ and XMTFSIZ must be programmed with the same value.
2. This field is write protected when SIU transmit is enabled.

7.2.8 Receive Clock Source Register

The receive clock source register is shown in Figure 7-8 and described in Table 7-22.

Figure 7-8   Receive Clock Source Register (0x000000C0)

31	1	0
Reserved		RCVSRC
R-0x00000000		R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-22   Receive Clock Source Register Field Descriptions

Bit	Field	Value	Description
31-1	Reserved	0x00000000	Reserved
0	RCVSRC		Receive clock and frame sync selection
		0	TCLK_A and TFS_A
		1	TCLK_B and TFS_B

7.2.9 Receive Control Register

The receive control register is shown in Figure 7-9 and described in Table 7-23.

Figure 7-9   Receive Control Register (0x000000C4)

31	30	29						16	15			8
Reserved		RCVDATD							Reserved			
R-0x0		R/W-0x0000							R-0x00			

7	6	5	4		3	2	1		0
RCVFSYNCP	RCVFCLKP	RCVDCLKP	Reserved				RCVDATR		RCVCLKM
R/W-0	R/W-0	R/W-0	R-00				R/W-00		R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-23   Receive Control Register Field Descriptions([1])

Bit	Field	Value	Description
31-30	Reserved	0x0	Reserved
29-16	RCVDATD([2])	0x0000	Receive data delay [0-16383]
			n + 1.5 When clock polarity for data and frame sync differs.
			n + 2 When clock polarity for data and frame sync is identical.
15 -8	Reserved	0x00	Reserved
7	RCVFSYNCP([2])		Receive frame sync polarity
		0	Receive frame sync is active low.
		1	Receive frame sync is active high.
6	RCVFCLKP([2])		Receive frame sync clock polarity
		0	Receive frame sync is sampled on the rising edge of TCLK_x.
		1	Receive frame sync is sampled on the falling edge of TCLK_x.
5	RCVDCLKP([2])		Receive data clock polarity
		0	Receive data is sampled on the rising edge of TCLK_x.
		1	Receive data is sampled on the falling edge of TCLK_x.
4-3	Reserved	0	Reserved
2-1	RCVDATR([2])		Receive data rate
		00	8.192 Mbps
		01	16.384 Mbps
		10	32.768 Mbps
		11	Reserved
0	RCVCLKM([2])		Receive clock mode
		0	TCLK_x is a double rate clock.
		1	TCLK_x is a single rate clock.

1. If CLKD is programmed to 0 or RCVSRC and XMTSRC will be written with the same value, then RCVFSYNCP and XMTFSYNCP must be programmed with the same value. Likewise, RCVFCLKP and XMTFCLKP must be programmed with the same value.

2. This field is write protected when SIU receive is enabled.

7.2.10 Receive Size Register

The receive size register is shown in Figure 7-10 and described in Table 7-24.

Figure 7-10   Receive Size Register (0x000000C8)

31	24	23	16	15	7	6	0
Reserved		RCVFCNT		Reserved		RCVFSIZ	
R-0x000		R/W-0x3		R-0x00		R/W-0x07F	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-24   Receive Size Register Field Descriptions[1]

Bit	Field	Value	Description
31-24	Reserved	0x000	Reserved
23-16	RCVFCNT[2]	0x3	Receive frame count [0-255]
			n + 1 Number of frames per superframe
15-7	Reserved	0x00	Reserved
6-0	RCVFSIZ[2]	0x7F	Receive frame size [0-27]
			n + 1 Number of logical timeslots per frame. Minimum number of logical timeslot periods per frame.
			≤6 Values of 6 and less are reserved. Attempts to write a value of 6 or less will result in a value of 0x7F being registered.

1. If CLKD is programmed to 0 or RCVSRC and XMTSRC will be written with the same value, then RCVFSIZ and XMTFSIZ must be programmed with the same value.

2. This field is write protected when SIU receive is enabled.

7.2.11 TDMU Global Control Register

The TDMU global control register is shown in Figure 7-11 and described in Table 7-25.

Figure 7-11   TDMU Global Control Register (0x00000100)

31	1	0
Reserved		DMAENB
R-0x00000000		R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-25   TDMU Global Control Register Field Descriptions

Bit	Field	Value	Description
31-1	Reserved		Reserved
0	DMAENB		TDMU and DMATCU enable
		0	TDMU and DMATCU is disabled/stopped.
		1	TDMU and DMATCU is enabled/running.

7.2.12 Transmit Free Running Frame Counter

The transmit free running frame counter is shown in Figure 7-12 and described in Table 7-26.

Figure 7-12   Transmit Free Running Frame Counter (0x00000104)

31	24	23	0
Reserved		FRFC	
R-0x00		R-0x000000	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-26   Transmit Free Running Frame Counter Field Descriptions

Bit	Field	Value	Description
31-24	Reserved	0x00	Reserved
23-0	FRFC	0x000000	Free running frame counter

7.2.13 Receive Free Running Frame Counter

The receive free running frame counter is shown in Figure 7-13 and described in Table 7-27.

Figure 7-13   Receive Free Running Frame Counter (0x00000108)

31	24	23	0
Reserved		FRFC	
R-0x00		R-0x000000	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-27   Receive Free Running Frame Counter Field Descriptions

Bit	Field	Value	Description
31-24	Reserved	0x00	Reserved
23-0	FRFC	0x000000	Free running frame counter

7.2.14 TDMU Global Configuration Register

The TDMU global configuration register is shown in Figure 7-14 and described in Table 7-28.

Figure 7-14   TDMU Global Configuration Register (0x0000010C)

31	1	0
Reserved		BEND
R-0x00000000		R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset
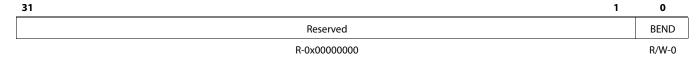
Table 7-28   TDMU Global Configuration Register Field Descriptions

Bit	Field	Value	Description
31-1	Reserved	0x00000000	Reserved
0	BEND(1)		Endian mode selection
		0	Little endian format
		1	Big endian format

1. This field is write protected once DMA is enabled.

7.2.15 Transmit Channel Bitmap Active Status Register

The transmit channel bitmap active status register is shown in Figure 7-15 and described in Table 7-29.

Figure 7-15   Transmit Channel Bitmap Active Status Register (0x00000110)
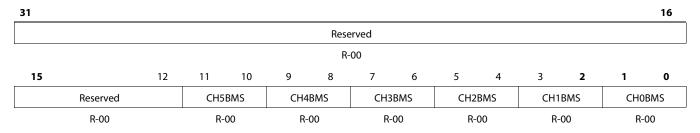
31														16
Reserved														
R-00														

15		12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved			CH5BMS		CH4BMS		CH3BMS		CH2BMS		CH1BMS		CH0BMS	
R-00			R-00		R-00		R-00		R-00		R-00		R-00	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-29   Transmit Channel Bitmap Active Status Register Field Descriptions (Part 1 of 2)

Bit	Field	Value	Description
31-12	Reserved	0x00000	Reserved
11-10	CH5BMS		Transmit channel 5 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved
9-8	CH4BMS		Transmit channel 4 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved
7-6	CH3BMS		Transmit channel 3 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved

Table 7-29   Transmit Channel Bitmap Active Status Register Field Descriptions (Part 2 of 2)
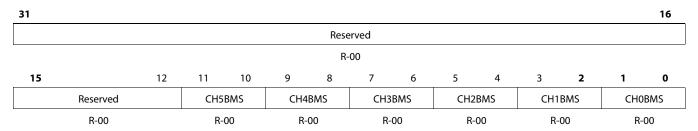
Bit	Field	Value	Description
5-4	CH2BMS		Transmit channel 2 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved
3-2	CH1BMS		Transmit channel 1 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved
1-0	CH0BMS		Transmit channel 0 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active
		11	Reserved

7.2.16 Receive Channel Bitmap Active Status Register

The receive channel bitmap active status register is shown in Figure 7-16 and described in Table 7-30.

Figure 7-16   Receive Channel Bitmap Active Status Register (0x00000114)

31															16
Reserved															
R-00															

15			12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved				CH5BMS		CH4BMS		CH3BMS		CH2BMS		CH1BMS		CH0BMS	
R-00				R-00		R-00		R-00		R-00		R-00		R-00	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-30   Receive Channel Bitmap Active Status Register Field Descriptions (Part 1 of 2)

Bit	Field	Value	Description
31-12	Reserved	0x00000	Reserved
11-10	CH5BMS		Receive channel 5 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved
9-8	CH4BMS		Receive channel 4 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved

Table 7-30   Receive Channel Bitmap Active Status Register Field Descriptions (Part 2 of 2)

Bit	Field	Value	Description
7-6	CH3BMS		Receive channel 3 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved
5-4	CH2BMS		Receive channel 2 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved
3-2	CH1BMS		Receive channel 1 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active.
		11	Reserved
1-0	CH0BMS		Receive channel 0 active status. Defines whether A and B bitmap is active or inactive.
		00	Both A and B bitmaps are inactive.
		01	A bitmap is active. B bitmap is inactive.
		10	A bitmap is inactive. B bitmap is active
		11	Reserved

7.2.17 DMATCU Global Control Register

The DMATCU global control register is shown in Figure 7-17 and described in Table 7-31.

**Figure 7-17   DMATCU Global Control Register (0x00000180)**

31						16
Reserved						
R-0x000000						

15	7	6	4	3	2	0
Reserved		MAXPRI		Rsvd	PRI	
R-0x000000		R/W-0x7		R-0b	R/W-0x7	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset
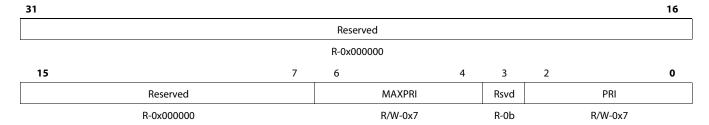
Table 7-31   DMATCU Global Control Register Field Descriptions

Bit	Field	Value	Description
31-7	Reserved	0x000000	Reserved
6-4	MAXPRI[1]		Maximum Transfer Priority. MAXPRI is expected to be programmed with a value less than or equal to PRI. When MAXPRI is greater than PRI, it is treated as if it were equal and no escalation or de-escalation takes place.
		0	Priority 0-Highest priority
		1	Priority 1
		...	Priority 2-6
		7	Priority 7-Lowest priority.
3	Reserved	0b	Reserved
2-0	PRI[1]		Transfer Priority.
		0	Priority 0-Highest priority
		1	Priority 1
		...	Priority 2-6
		7	Priority 7-Lowest priority

1. This field is write protected once DMA is enabled.

7.2.18 Transmit Timeslot Delay Register

The transmit timeslot delay register is shown in Figure 7-18 and described in Table 7-32.

Figure 7-18   Transmit Timeslot Delay Register (0x00000184)

31							16
Reserved							
R-0x000000							

15	14	13	12	11	7	6	0
XMTSFINT		XMTFRINT		Reserved		XMTFDLY	
R-00b		R-00b		R-00000b		R/W-0x07F	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-32   Transmit Timeslot Register Field Descriptions

Bit	Field	Value	Description
31-16	Reserved	0x000000	Reserved
15-14	XMTSFINT[1]		Transmit superframe interrupt selection
		00	Interrupt asserted on transfer acknowledge only.
		01	Interrupt asserted on delay count only.
		10	Interrupt asserted on either transfer acknowledge or delay count.
		11	Interrupt asserted on both transfer acknowledge and delay count.

Table 7-32   Transmit Timeslot Register Field Descriptions

Bit	Field	Value	Description
13-12	XMTFRINT([1])		Transmit frame interrupt selection
		00	Interrupt asserted on transfer acknowledge only.
		01	Interrupt asserted on delay count only.
		10	Interrupt asserted on either transfer acknowledge or delay count.
		11	Interrupt asserted on both transfer acknowledge and delay count.
11-7	Reserved	00000b	Reserved
6-0	XMTFDLY([1])	0x07F	Transfer frame interrupt delay [0-127]
			n + 1 Number of logical timeslots the frame interrupt is delayed after going pending.

1. This field is write protected once DMA is enabled.

7.2.19 Receive Timeslot Delay Register

The receive timeslot delay register is shown in Figure 7-19 and described in Table 7-33.

Figure 7-19   Receive Timeslot Delay Register (0x00000188)

31							16
				Reserved			
				R-0x000000			

15	14	13	12	11	7	6	0
RCVSFINT		RCVFRINT		Reserved		RCVFDLY	
R-00b		R-00b		R-00000b		R/W-0x07F	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-33   Receive Timeslot Delay Register Field Descriptions

Bit	Field	Value	Description
31-16	Reserved	0x000000	Reserved
15-14	RCVSFINT(1)		Receive superframe interrupt selection
		00	Interrupt asserted on transfer acknowledge only.
		01	Interrupt asserted on delay count only.
		10	Interrupt asserted on either transfer acknowledge or delay count.
		11	Interrupt asserted on both transfer acknowledge and delay count.
13-12	RCVFRINT(1)		Receive frame interrupt selection
		00	Interrupt asserted on transfer acknowledge only.
		01	Interrupt asserted on delay count only.
		10	Interrupt asserted on either transfer acknowledge or delay count.
		11	Interrupt asserted on both transfer acknowledge and delay count.
11-7	Reserved	00000b	Reserved
6-0	RCVFDLY(1)	0x07F	Receive frame interrupt delay [0-127]
			n + 1 Number of logical timeslots the frame interrupt is delayed after going pending.

1. This field is write protected once DMA is enabled.

7.2.20 Transmit Channel Configuration Active Status Register

The transmit channel configuration active status register is shown in Figure 7-20 and described in Table 7-34.

Figure 7-20   Transmit Channel Configuration Active Status Register (0x00000190)

31		16
	Reserved	

R-00000

15	12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved		CH5CS		CH4CS		CH3CS		CH2CS		CH1CS		CH0CS	
R-00000		R-00		R-00		R-00		R-00		R-00		R-00	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-34   Transmit Channel Configuration Active Status Register Field Descriptions (Part 1 of 2)

Bit	Field	Value	Description
31-12	Reserved	0x00000	Reserved
11-10	CH5CS		Transmit channel 5 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved
9-8	CH4CS		Transmit channel 4 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved
7-6	CH3CS		Transmit channel 3 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved

Table 7-34   Transmit Channel Configuration Active Status Register Field Descriptions (Part 2 of 2)

Bit	Field	Value	Description
5-4	CH2CS		Transmit channel 2 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved
3-2	CH1CS		Transmit channel 1 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved
1-0	CH0CS		Transmit channel 0 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved

7.2.21 Receive Channel Configuration Active Status Register

The receive channel configuration active status register is shown in Figure 7-21 and described in Table 7-35.

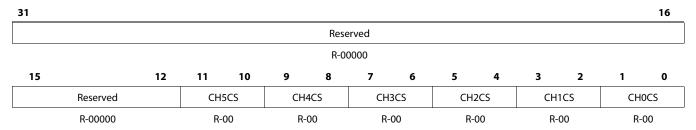**Figure 7-21   Receive Channel Configuration Active Status Register (0x00000194)**

31															16
Reserved															
R-00000															

15			12	11	10	9	8	7	6	5	4	3	2	1	0
Reserved				CH5CS		CH4CS		CH3CS		CH2CS		CH1CS		CH0CS	
R-00000				R-00		R-00		R-00		R-00		R-00		R-00	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-35   Receive Channel Configuration Active Status Register Field Descriptions (Part 1 of 2)

Bit	Field	Value	Description
31-12	Reserved	0x00000	Reserved
11-10	CH5CS		Receive channel 5 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved
9-8	CH4CS		Receive channel 4 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved

Table 7-35   Receive Channel Configuration Active Status Register Field Descriptions (Part 2 of 2)

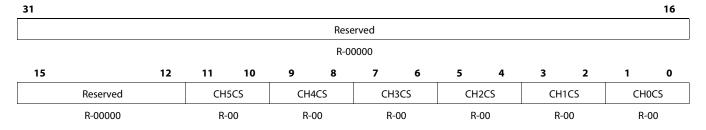| Bit	Field	Value	Description
7-6	CH3CS		Receive channel 3 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved
5-4	CH2CS		Receive channel 2 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved
3-2	CH1CS		Receive channel 1 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved
1-0	CH0CS		Receive channel 0 active status. Defines whether the A and B configurations are active or inactive.
		00	Both A and B configurations are inactive.
		01	A configuration is active. B configuration is inactive.
		10	A configuration is inactive. B configuration is active.
		11	Reserved

7.2.22 Channel n Error Control Register

The channel n error control register is shown in Figure 7-22 and described in Table 7-36.

Figure 7-22   Channel *n* Error Control Register

31		16
Reserved		
R-0x000000		

15	9	8	7	2	1	0
Reserved		CEQOV	Reserved		CLRQ	POPQ
R-0x000000		R/W-0	R-0x00		R/W-0	R/W-0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-36   Channel *n* Error Control Register Field Descriptions

Bit	Field	Value	Description
31-9	Reserved	0x000000	Reserved
8	CEQOV		Clear error queue overflow flag. This bit is always read as 0.
		0	Write 0 has no effect.
		1	Write 1 to clear the error queue overflow flag in the error count register.
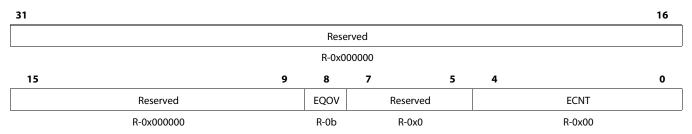
Table 7-36   Channel *n* Error Control Register Field Descriptions

Bit	Field	Value	Description
7-2	Reserved	0x00	Reserved
1	CLRQ		Clear error queue. This bit is always read as 0.
		0	Write 0 has no effect.
		1	Write 1 to clear the error queue of all error entries.
0	POPQ		Pop the top entry off the queue. This bit is always read as 0.
		0	Write 0 has no effect.
		1	Write 1 to pop the top error entry off the error queue.

7.2.23 Channel *n* Error Count Register

The channel *n* error count register is shown in Figure 7-23 and described in Table 7-37.
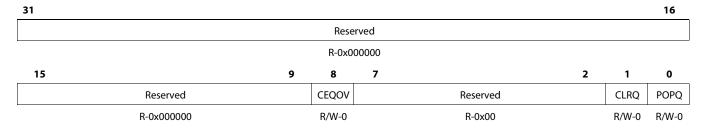
Figure 7-23   Channel *n* Error Count Register

31						16
Reserved						
R-0x000000						

15		9	8	7	5	4	0
Reserved			EQOV	Reserved		ECNT	
R-0x000000			R-0b	R-0x0		R-0x00	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset
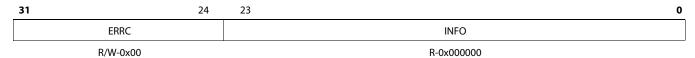
Table 7-37   Channel *n* Error Count Register Field Descriptions

Bit	Field	Value	Description
31-9	Reserved	0x000000	Reserved
8	EQOV		Error Queue Overflow. The overflow flag is set to indicate that the error detection logic was unable to add at least one error entry to the error queue. Once the overflow flag is set it remains set until it is cleared by a write to the overflow clear bit in the error queue control register. Reading entries from the queue does not automatically clear the overflow flag.
7-5	Reserved	0x0	Reserved
4-0	ECNT		Error Count, the number of unread entries on the error queue.

7.2.24 Channel *n* Error Queue Register

The channel *n* error queue register is shown in Figure 7-24 and described in Table 7-38.

Figure 7-24   Channel *n* Error Queue Register

31	24	23	0
ERRC		INFO	
R/W-0x00		R-0x000000	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-38   Channel *n* Error Queue Register Field Descriptions

Bit	Field	Value	Description
31-24	ERRC		Error Code
23-0	INFO		Error Information Field. Information specific to the cause of the error is provided in this field. The specific format is error specific as defined by the following error code figures.

7.2.25 CAT-0 Error Code Format

The CAT-0 error code format is shown in Figure 7-25 and described in Table 7-39.

Figure 7-25   CAT-0 Error Code Format

31	24	23	0
ERRC		Reserved	
R/W-0x00		R/W-0x000000	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-39   CAT-0 Error Code Format Field Descriptions

Bit	Field	Value	Description
31-24	ERRC		Error Code
23-0	Reserved	0x000000	Reserved

7.2.26 CAT-1 Error Code Format

The CAT-1 error code format is shown in Figure 7-26 and described in Table 7-40.

Figure 7-26   CAT-1 Error Code Format

31	24	23	0
ERRC		FUNM	
R/W-0x00		R/W-0x000000	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-40   CAT-1 Error Code Format Field Descriptions

Bit	Field	Value	Description
31-24	ERRC	0x00	Error code
23-0	FNUM	0x000000	Free running frame counter value

7.2.27 CAT-2 Error Code Format

The CAT-2 error code format is shown in Figure 7-27 and described in Table 7-41.

Figure 7-27   CAT-2 Error Code Format

31		24	23		16	15		8	7		0
ERRC			Reserved			CIDA			CIDF		
R/W-0x00			R/W-0x00			R/W-0x000			R/W-0x00		

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset
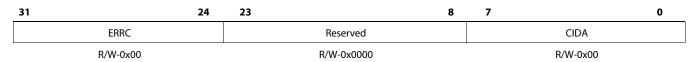
Table 7-41   CAT-2 Error Code Format Field Descriptions

Bit	Field	Value	Description
31-24	ERRC	0x00	Error code
23-16	Reserved	0x00	Reserved
15-8	CIDA	0x00	Active configuration ID
7-0	CIDF	0x00	Data buffer configuration ID

7.2.28 CAT-3 Error Code Format

The CAT-3 error code format is shown in Figure 7-28 and described in Table 7-42.

Figure 7-28   CAT-3 Error Code Format

31		24	23		8	7		0
ERRC			Reserved			CIDA		
R/W-0x00			R/W-0x0000			R/W-0x00		

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-42   CAT-3 Error Code Format Field Descriptions

Bit	Field	Value	Description
31-24	ERRC	0x00	Error code
23-8	Reserved	0x0000	Reserved
7-0	CIDA	0x00	Active configuration ID

7.2.29 Channel *n* Enable Register

The channel *n* enable register is shown in Figure 7-29 and described in Table 7-43.

Figure 7-29   Channel *n* Enable Register

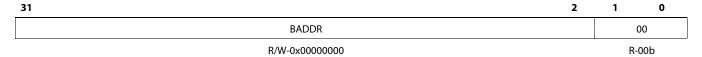| 31		16	15		8	7		1	0
Reserved			CnID			Reserved			CHnENB
R-0x0000			R-0x00			R-0x00			R-0x0

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-43   Channel *n* Enable Register Field Descriptions

Bit	Field	Value	Description
31-16	Reserved	0x0000	Reserved
16-8	CnID		Configuration Identification. Defines the configuration selection that takes effect on the next superframe synchronization event. The least significant bit determines whether the A or the B channel bitmap and configuration registers are used.
7-1	Reserved	0x00b	Reserved
0	CHnENB		Channel n enable.
		0	Channel n is disabled/stopped.
		1	Channel n is enabled/running.

7.2.30 Channel *n* Memory Address Register

The channel *n* memory address register is shown in Figure 7-30 and described in Table 7-44.

Figure 7-30   Channel *n* Memory Address Register

31	2	1	0
BADDR		00	
R/W-0x00000000		R-00b	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-44   Channel *n* Memory Address Register Field Descriptions

Bit	Field	Value	Description
31-0	BADDR	0x00000000	Memory Buffer Base Address. The two least significant bits are always forced to zero.

7.2.31 Channel *n* Frame Allocation Register

The channel *n* frame allocation register is shown in Figure 7-31 and described in Table 7-45.
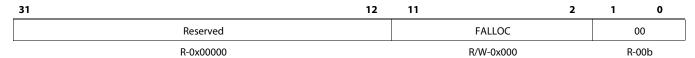
Figure 7-31   Channel *n* Frame Allocation Register

31	12	11	2	1	0
Reserved		FALLOC		00	
R-0x00000		R/W-0x000		R-00b	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-45   Channel *n* Frame Allocation Register Field Descriptions
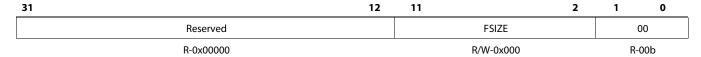
Bit	Field	Value	Description
31-12	Reserved	0x00000	Reserved
11-0	FALLOC		Frame buffer allocation. Each frame in the memory buffer begins on successive increments of FALLOC from BADDR. The two least significant bits are always forced to zero.

7.2.32 Channel *n* Frame Size Register

The channel *n* frame size register is shown in Figure 7-32 and described in Table 7-46.

Figure 7-32   Channel *n* Frame Size Register

31		12	11		2	1	0
Reserved			FSIZE			00	
R-0x00000			R/W-0x000			R-00b	

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset
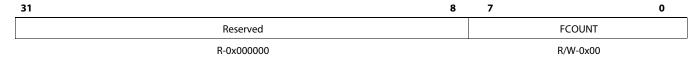
Table 7-46   Channel *n* Frame Size Register Field Descriptions

Bit	Field	Value	Description
31-12	Reserved	0x00000	Reserved
11-0	FSIZE		Frame size. Number of bytes to transfer for each frame. Applies only to transmit channels; ignored for receive channels. The two least significant bits are always forced to zero.

7.2.33 Channel *n* Frame Count Register

The channel *n* frame count register is shown in Figure 7-33 and described in Table 7-47.

Figure 7-33   Channel *n* Frame Count Register

31		8	7		0
Reserved			FCOUNT		
R-0x000000			R/W-0x00		

LEGEND: R/W = Read/Write; R = Read only; -n = value after reset

Table 7-47   Channel *n* Frame Count Register Field Descriptions

Bit	Field	Value	Description
31-8	Reserved	0x000000	Reserved
7-0	FCOUNT		Frame Count. Number of frames in a buffer. Appropriate values are 1 through 255. Every FCOUNT frames the location of the frame in the memory buffer wraps back to BADDR.

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